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*****
174341 Thu Sep 1 02:19:26 2011
new/src/share/vm/opto/memnode.cpp
*****
_____unchanged_portion_omitted_____

1485 //-----Value-----
1486 const Type *LoadNode::Value( PhaseTransform *phase ) const {
1487 // Either input is TOP ==> the result is TOP
1488 Node* mem = in(MemNode::Memory);
1489 const Type *t1 = phase->type(mem);
1490 if (t1 == Type::TOP) return Type::TOP;
1491 Node* adr = in(MemNode::Address);
1492 const TypePtr* tp = phase->type(adr)->isa_ptr();
1493 if (tp == NULL || tp->empty()) return Type::TOP;
1494 int off = tp->offset();
1495 assert(off != Type::OffsetTop, "case covered by TypePtr::empty");
1496 Compile* C = phase->C;
1497 #endif /* ! codereview */

1499 // Try to guess loaded type from pointer type
1500 if (tp->base() == Type::AryPtr) {
1501     const Type *t = tp->is_aryptr()->elem();
1502     // Don't do this for integer types. There is only potential profit if
1503     // the element type t is lower than _type; that is, for int types, if _type
1504     // more restrictive than t. This only happens here if one is short and the
1505     // char (both 16 bits), and in those cases we've made an intentional decisio
1506     // to use one kind of load over the other. See AndINode::Ideal and 4965907.
1507     // Also, do not try to narrow the type for a LoadKlass, regardless of offset
1508     //
1509     // Yes, it is possible to encounter an expression like (LoadKlass p1:(AddP x
1510     // where the _gvn.type of the AddP is wider than 8. This occurs when an ear
1511     // copy p0 of (AddP x x 8) has been proven equal to p1, and the p0 has been
1512     // subsumed by p1. If p1 is on the worklist but has not yet been re-transfo
1513     // it is possible that p1 will have a type like Foo*[int+]:NotNull*+any.
1514     // In fact, that could have been the original type of p1, and p1 could have
1515     // had an original form like p1:(AddP x x (LShiftL quux 3)), where the
1516     // expression (LShiftL quux 3) independently optimized to the constant 8.
1517     if ((t->isa_int() == NULL) && (t->isa_long() == NULL)
1518         && Opcode() != Op_LoadKlass && Opcode() != Op_LoadNklass) {
1519         // t might actually be lower than _type, if _type is a unique
1520         // concrete subclass of abstract class t.
1521         // Make sure the reference is not into the header, by comparing
1522         // the offset against the offset of the start of the array's data.
1523         // Different array types begin at slightly different offsets (12 vs. 16).
1524         // We choose T_BYTE as an example base type that is least restrictive
1525         // as to alignment, which will therefore produce the smallest
1526         // possible base offset.
1527         const int min_base_off = arrayOopDesc::base_offset_in_bytes(T_BYTE);
1528         if ((uint)off >= (uint)min_base_off) { // is the offset beyond the header
1529             const Type* jt = t->join(_type);
1530             // In any case, do not allow the join, per se, to empty out the type.
1531             if (jt->empty() && !t->empty()) {
1532                 // This can happen if a interface-typed array narrows to a class type.
1533                 jt = _type;
1534             }
1535         }
1536     }
1537     if (EliminateAutoBox && adr->is_AddP()) {
1538         // The pointers in the autobox arrays are always non-null
1539         Node* base = adr->in(AddPNode::Base);
1540         if (base != NULL &&
1541             !phase->type(base)->higher_equal(TypePtr::NULL_PTR)) {
1542             Compile::AliasType* atp = C->alias_type(base->adr_type());
1543             Compile::AliasType* atp = phase->C->alias_type(base->adr_type());
1544             if (is_autobox_cache(atp)) {
1545                 return jt->join(TypePtr::NOTNULL)->is_ptr();

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1544     }
1545     }
1546     }
1547     return jt;
1548 }
1549 }
1550 } else if (tp->base() == Type::InstPtr) {
1551     ciEnv* env = C->env();
1552     #endif /* ! codereview */
1553     const TypeInstPtr* tinst = tp->is_instptr();
1554     ciKlass* klass = tinst->klass();
1555     assert( off != Type::OffsetBot ||
1556            // arrays can be cast to Objects
1557            tp->is_oopptr()->klass()->is_java_lang_Object() ||
1558            // unsafe field access may not have a constant offset
1559            C->has_unsafe_access(),
1560            "Field accesses must be precise" );
1561     // For oop loads, we expect the _type to be precise
1562     if (klass == env->String_klass() &&
1563         if (klass == phase->C->env()->String_klass() &&
1564             adr->is_AddP() && off != Type::OffsetBot) {
1565         // For constant Strings treat the final fields as compile time constants.
1566         Node* base = adr->in(AddPNode::Base);
1567         const TypeOopPtr* t = phase->type(base)->isa_oopptr();
1568         if (t != NULL && t->singleton()) {
1569             ciField* field = env->String_klass()->get_field_by_offset(off, false);
1570             ciField* field = phase->C->env()->String_klass()->get_field_by_offset(of
1571             if (field != NULL && field->is_final()) {
1572                 ciObject* string = t->const_oop();
1573                 ciConstant constant = string->as_instance()->field_value(field);
1574                 if (constant.basic_type() == T_INT) {
1575                     return TypeInt::make(constant.as_int());
1576                 } else if (constant.basic_type() == T_ARRAY) {
1577                     if (adr->bottom_type()->is_ptr_to_narrowoop()) {
1578                         return TypeNarrowOop::make_from_constant(constant.as_object(), tru
1579                     } else {
1580                         return TypeOopPtr::make_from_constant(constant.as_object(), true);
1581                     }
1582                 }
1583             }
1584         }
1585     }
1586     // Optimizations for constant objects
1587     ciObject* const_oop = tinst->const_oop();
1588     if (const_oop != NULL) {
1589         // For constant CallSites treat the target field as a compile time constan
1590         if (const_oop->is_call_site()) {
1591             ciCallSite* call_site = const_oop->as_call_site();
1592             ciField* field = call_site->klass()->as_instance_klass()->get_field_by_o
1593             if (field != NULL && field->is_call_site_target()) {
1594                 ciMethodHandle* target = call_site->get_target();
1595                 if (target != NULL) { // just in case
1596                     ciConstant constant(T_OBJECT, target);
1597                     const Type* t;
1598                     if (adr->bottom_type()->is_ptr_to_narrowoop()) {
1599                         t = TypeNarrowOop::make_from_constant(constant.as_object(), true);
1600                     } else {
1601                         t = TypeOopPtr::make_from_constant(constant.as_object(), true);
1602                     }
1603                 }
1604                 // Add a dependence for invalidation of the optimization.
1605                 if (!call_site->is_constant_call_site()) {
1606                     C->dependencies()->assert_call_site_target_value(call_site, target
1607                 }
1608                 return t;
1609             }
1610         }

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1607     }
1608   }
1609 }
1610 #endif /* ! codereview */
1611 } else if (tp->base() == Type::KlassPtr) {
1612   assert( off != Type::OffsetBot ||
1613     // arrays can be cast to Objects
1614     tp->is_klassptr()->klass()->is_java_lang_Object() ||
1615     // also allow array-loading from the primary supertype
1616     // array during subtype checks
1617     Opcode() == Op_LoadKlass,
1618     "Field accesses must be precise" );
1619   // For klass/static loads, we expect the _type to be precise
1620 }

1622 const TypeKlassPtr *tkls = tp->isa_klassptr();
1623 if (tkls != NULL && !StressReflectiveCode) {
1624   ciKlass* klass = tkls->klass();
1625   if (klass->is_loaded() && tkls->klass_is_exact()) {
1626     // We are loading a field from a Klass metaobject whose identity
1627     // is known at compile time (the type is "exact" or "precise").
1628     // Check for fields we know are maintained as constants by the VM.
1629     if (tkls->offset() == Klass::super_check_offset_in_bytes() + (int)s
1630       // The field is Klass::_super_check_offset. Return its (constant) value
1631       // (Folds up type checking code.)
1632       assert(Opcode() == Op_LoadI, "must load an int from _super_check_offset"
1633         return TypeInt::make(klass->super_check_offset());
1634     }
1635     // Compute index into primary_supers array
1636     jint depth = (tkls->offset() - (Klass::primary_supers_offset_in_bytes() +
1637     // Check for overflowing; use unsigned compare to handle the negative case
1638     if( depth < ciKlass::primary_super_limit() ) {
1639       // The field is an element of Klass::_primary_supers. Return its (const
1640       // (Folds up type checking code.)
1641       assert(Opcode() == Op_LoadKlass, "must load a klass from _primary_supers
1642         ciKlass *ss = klass->super_of_depth(depth);
1643         return ss ? TypeKlassPtr::make(ss) : TypePtr::NULL_PTR;
1644     }
1645     const Type* aift = load_array_final_field(tkls, klass);
1646     if (aift != NULL) return aift;
1647     if (tkls->offset() == in_bytes(arrayKlass::component_mirror_offset()) + (i
1648       && klass->is_array_klass()) {
1649       // The field is arrayKlass::_component_mirror. Return its (constant) va
1650       // (Folds up aClassConstant.getComponentType, common in Arrays.copyOf.)
1651       assert(Opcode() == Op_LoadP, "must load an oop from _component_mirror");
1652       return TypeInstPtr::make(klass->as_array_klass()->component_mirror());
1653     }
1654     if (tkls->offset() == Klass::java_mirror_offset_in_bytes() + (int)sizeof(o
1655       // The field is Klass::_java_mirror. Return its (constant) value.
1656       // (Folds up the 2nd indirection in anObjConstant.getClass().)
1657       assert(Opcode() == Op_LoadP, "must load an oop from _java_mirror");
1658       return TypeInstPtr::make(klass->java_mirror());
1659     }
1660   }

1662 // We can still check if we are loading from the primary_supers array at a
1663 // shallow enough depth. Even though the klass is not exact, entries less
1664 // than or equal to its super depth are correct.
1665 if (klass->is_loaded()) {
1666   ciType *inner = klass->klass();
1667   while( inner->is_obj_array_klass() )
1668     inner = inner->as_obj_array_klass()->base_element_type();
1669   if( inner->is_instance_klass() &&
1670     !inner->as_instance_klass()->flags().is_interface() ) {
1671     // Compute index into primary_supers array
1672     jint depth = (tkls->offset() - (Klass::primary_supers_offset_in_bytes()

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1673 // Check for overflowing; use unsigned compare to handle the negative ca
1674 if( depth < ciKlass::primary_super_limit() &&
1675   depth <= klass->super_depth() ) { // allow self-depth checks to hand
1676   // The field is an element of Klass::_primary_supers. Return its (con
1677   // (Folds up type checking code.)
1678   assert(Opcode() == Op_LoadKlass, "must load a klass from _primary_supe
1679     ciKlass *ss = klass->super_of_depth(depth);
1680     return ss ? TypeKlassPtr::make(ss) : TypePtr::NULL_PTR;
1681   }
1682 }
1683 }

1685 // If the type is enough to determine that the thing is not an array,
1686 // we can give the layout_helper a positive interval type.
1687 // This will help short-circuit some reflective code.
1688 if (tkls->offset() == Klass::layout_helper_offset_in_bytes() + (int)sizeof(o
1689   && !klass->is_array_klass() // not directly typed as an array
1690   && !klass->is_interface() // specifically not Serializable & Cloneable
1691   && !klass->is_java_lang_Object() // not the supertype of all T[]
1692   ) {
1693   // Note: When interfaces are reliable, we can narrow the interface
1694   // test to (klass != Serializable && klass != Cloneable).
1695   assert(Opcode() == Op_LoadI, "must load an int from _layout_helper");
1696   jint min_size = Klass::instance_layout_helper(oopDesc::header_size(), fals
1697   // The key property of this type is that it folds up tests
1698   // for array-ness, since it proves that the layout_helper is positive.
1699   // Thus, a generic value like the basic object layout helper works fine.
1700   return TypeInt::make(min_size, max_jint, Type::WidenMin);
1701 }
1702 }

1704 // If we are loading from a freshly-allocated object, produce a zero,
1705 // if the load is provably beyond the header of the object.
1706 // (Also allow a variable load from a fresh array to produce zero.)
1707 const TypeOopPtr *tinst = tp->isa_copptr();
1708 bool is_instance = (tinst != NULL) && tinst->is_known_instance_field();
1709 if (ReduceFieldZeroing || is_instance) {
1710   Node* value = can_see_stored_value(mem, phase);
1711   if (value != NULL && value->is_Con())
1712     return value->bottom_type();
1713 }

1715 if (is_instance) {
1716   // If we have an instance type and our memory input is the
1717   // program's initial memory state, there is no matching store,
1718   // so just return a zero of the appropriate type
1719   Node *mem = in(MemNode::Memory);
1720   if (mem->is_Parm() && mem->in(0)->is_Start()) {
1721     assert(mem->as_Parm()->_con == TypeFunc::Memory, "must be memory Parm");
1722     return Type::get_zero_type(_type->basic_type());
1723   }
1724 }
1725 return _type;
1726 }

1728 //-----match_edge-----
1729 // Do we Match on this edge index or not? Match only the address.
1730 uint LoadNode::match_edge(uint idx) const {
1731   return idx == MemNode::Address;
1732 }

1734 //-----LoadBNode::Ideal-----
1735 //
1736 // If the previous store is to the same address as this load,
1737 // and the value stored was larger than a byte, replace this load
1738 // with the value stored truncated to a byte. If no truncation is

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1739 // needed, the replacement is done in LoadNode::Identity().
1740 //
1741 Node *LoadENode::Ideal(PhaseGVN *phase, bool can_reshape) {
1742     Node* mem = in(MemNode::Memory);
1743     Node* value = can_see_stored_value(mem, phase);
1744     if( value && !phase->type(value)->higher_equal(_type) ) {
1745         Node *result = phase->transform( new (phase->C, 3) LShiftINode(value, phase-
1746         return new (phase->C, 3) RShiftINode(result, phase->intcon(24));
1747     }
1748     // Identity call will handle the case where truncation is not needed.
1749     return LoadNode::Ideal(phase, can_reshape);
1750 }

1752 //-----LoadUBNode::Ideal-----
1753 //
1754 // If the previous store is to the same address as this load,
1755 // and the value stored was larger than a byte, replace this load
1756 // with the value stored truncated to a byte. If no truncation is
1757 // needed, the replacement is done in LoadNode::Identity().
1758 //
1759 Node* LoadUBNode::Ideal(PhaseGVN* phase, bool can_reshape) {
1760     Node* mem = in(MemNode::Memory);
1761     Node* value = can_see_stored_value(mem, phase);
1762     if( value && !phase->type(value)->higher_equal(_type) )
1763         return new (phase->C, 3) AndINode(value, phase->intcon(0xFF));
1764     // Identity call will handle the case where truncation is not needed.
1765     return LoadNode::Ideal(phase, can_reshape);
1766 }

1768 //-----LoadUSNode::Ideal-----
1769 //
1770 // If the previous store is to the same address as this load,
1771 // and the value stored was larger than a char, replace this load
1772 // with the value stored truncated to a char. If no truncation is
1773 // needed, the replacement is done in LoadNode::Identity().
1774 //
1775 Node *LoadUSNode::Ideal(PhaseGVN *phase, bool can_reshape) {
1776     Node* mem = in(MemNode::Memory);
1777     Node* value = can_see_stored_value(mem, phase);
1778     if( value && !phase->type(value)->higher_equal(_type) )
1779         return new (phase->C, 3) AndINode(value, phase->intcon(0xFFFF));
1780     // Identity call will handle the case where truncation is not needed.
1781     return LoadNode::Ideal(phase, can_reshape);
1782 }

1784 //-----LoadSNode::Ideal-----
1785 //
1786 // If the previous store is to the same address as this load,
1787 // and the value stored was larger than a short, replace this load
1788 // with the value stored truncated to a short. If no truncation is
1789 // needed, the replacement is done in LoadNode::Identity().
1790 //
1791 Node *LoadSNode::Ideal(PhaseGVN *phase, bool can_reshape) {
1792     Node* mem = in(MemNode::Memory);
1793     Node* value = can_see_stored_value(mem, phase);
1794     if( value && !phase->type(value)->higher_equal(_type) ) {
1795         Node *result = phase->transform( new (phase->C, 3) LShiftINode(value, phase-
1796         return new (phase->C, 3) RShiftINode(result, phase->intcon(16));
1797     }
1798     // Identity call will handle the case where truncation is not needed.
1799     return LoadNode::Ideal(phase, can_reshape);
1800 }

1802 //-----LoadKlassNode::make-----
1803 //
1804 // Polymorphic factory method:

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1805 Node *LoadKlassNode::make( PhaseGVN& gvn, Node *mem, Node *adr, const TypePtr* a
1806     Compile* C = gvn.C;
1807     Node *ctl = NULL;
1808     // sanity check the alias category against the created node type
1809     const TypeOopPtr *adr_type = adr->bottom_type()->isa_oopptr();
1810     assert(adr_type != NULL, "expecting TypeOopPtr");
1811     #ifdef_LP64
1812     if (adr_type->is_ptr_to_narrowoop()) {
1813         Node* load_klass = gvn.transform(new (C, 3) LoadNKlassNode(ctl, mem, adr, at
1814         return new (C, 2) DecodeNNode(load_klass, load_klass->bottom_type()->make_pt
1815     }
1816     #endif
1817     assert(!adr_type->is_ptr_to_narrowoop(), "should have got back a narrow oop");
1818     return new (C, 3) LoadKlassNode(ctl, mem, adr, at, tk);
1819 }

1821 //-----Value-----
1822 const Type *LoadKlassNode::Value( PhaseTransform *phase ) const {
1823     return klass_value_common(phase);
1824 }

1826 const Type *LoadNode::klass_value_common( PhaseTransform *phase ) const {
1827     // Either input is TOP ==> the result is TOP
1828     const Type *t1 = phase->type( in(MemNode::Memory) );
1829     if (t1 == Type::TOP) return Type::TOP;
1830     Node *adr = in(MemNode::Address);
1831     const Type *t2 = phase->type( adr );
1832     if (t2 == Type::TOP) return Type::TOP;
1833     const TypePtr *tp = t2->is_ptr();
1834     if (TypePtr::above_centerline(tp->ptr()) ||
1835         tp->ptr() == TypePtr::Null) return Type::TOP;

1837     // Return a more precise klass, if possible
1838     const TypeInstPtr *tinst = tp->isa_instptr();
1839     if (tinst != NULL) {
1840         ciInstanceKlass* ik = tinst->klass()->as_instance_klass();
1841         int offset = tinst->offset();
1842         if (ik == phase->C->env()->Class_klass()
1843             && (offset == java_lang_Class::klass_offset_in_bytes() ||
1844                 offset == java_lang_Class::array_klass_offset_in_bytes())) {
1845             // We are loading a special hidden field from a Class mirror object,
1846             // the field which points to the VM's Klass metaobject.
1847             ciType* t = tinst->java_mirror_type();
1848             // java_mirror_type returns non-null for compile-time Class constants.
1849             if (t != NULL) {
1850                 // constant oop => constant klass
1851                 if (offset == java_lang_Class::array_klass_offset_in_bytes()) {
1852                     return TypeKlassPtr::make(ciArrayKlass::make(t));
1853                 }
1854                 if (!t->is_klass()) {
1855                     // a primitive Class (e.g., int.class) has NULL for a klass field
1856                     return TypePtr::NULL_PTR;
1857                 }
1858                 // (Folds up the 1st indirection in aClassConstant.getModifiers().)
1859                 return TypeKlassPtr::make(t->as_klass());
1860             }
1861             // non-constant mirror, so we can't tell what's going on
1862         }
1863     }
1864     if( !ik->is_loaded() )
1865         return _type; // Bail out if not loaded
1866     if (offset == oopDesc::klass_offset_in_bytes()) {
1867         if (tinst->klass_is_exact()) {
1868             return TypeKlassPtr::make(ik);
1869         }
1870         // See if we can become precise: no subclasses and no interface
1871         // (Note: We need to support verified interfaces.)

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1871     if (!ik->is_interface() && !ik->has_subclass()) {
1872         //assert(!UseExactTypes, "this code should be useless with exact types")
1873         // Add a dependency; if any subclass added we need to recompile
1874         if (!ik->is_final()) {
1875             // %% should use stronger assert_unique_concrete_subtype instead
1876             phase->C->dependencies()->assert_leaf_type(ik);
1877         }
1878         // Return precise class
1879         return TypeKlassPtr::make(ik);
1880     }

1882     // Return root of possible class
1883     return TypeKlassPtr::make(TypePtr::NotNull, ik, 0/*offset*/);
1884 }
1885 }

1887 // Check for loading class from an array
1888 const TypeAryPtr *tary = tp->isa_aryptr();
1889 if( tary != NULL ) {
1890     ciKlass *tary_klass = tary->klass();
1891     if (tary_klass != NULL // can be NULL when at BOTTOM or TOP
1892         && tary->offset() == oopDesc::klass_offset_in_bytes() ) {
1893         if (tary->klass_is_exact()) {
1894             return TypeKlassPtr::make(tary_klass);
1895         }
1896         ciArrayKlass *ak = tary->klass()->as_array_klass();
1897         // If the klass is an object array, we defer the question to the
1898         // array component klass.
1899         if( ak->is_obj_array_klass() ) {
1900             assert( ak->is_loaded(), "" );
1901             ciKlass *base_k = ak->as_obj_array_klass()->base_element_klass();
1902             if( base_k->is_loaded() && base_k->is_instance_klass() ) {
1903                 ciInstanceKlass* ik = base_k->as_instance_klass();
1904                 // See if we can become precise: no subclasses and no interface
1905                 if (!ik->is_interface() && !ik->has_subclass()) {
1906                     //assert(!UseExactTypes, "this code should be useless with exact typ
1907                     // Add a dependency; if any subclass added we need to recompile
1908                     if (!ik->is_final()) {
1909                         phase->C->dependencies()->assert_leaf_type(ik);
1910                     }
1911                     // Return precise array class
1912                     return TypeKlassPtr::make(ak);
1913                 }
1914             }
1915             return TypeKlassPtr::make(TypePtr::NotNull, ak, 0/*offset*/);
1916         } else { // Found a type-array?
1917             //assert(!UseExactTypes, "this code should be useless with exact types")
1918             assert( ak->is_type_array_klass(), "" );
1919             return TypeKlassPtr::make(ak); // These are always precise
1920         }
1921     }
1922 }

1924 // Check for loading class from an array klass
1925 const TypeKlassPtr *tkls = tp->isa_klassptr();
1926 if (tkls != NULL && !StressReflectiveCode) {
1927     ciKlass* klass = tkls->klass();
1928     if( !klass->is_loaded() )
1929         return_type; // Bail out if not loaded
1930     if( klass->is_obj_array_klass() &&
1931         (uint)tkls->offset() == objArrayKlass::element_klass_offset_in_bytes() +
1932         ciKlass* elem = klass->as_obj_array_klass()->element_klass();
1933         // // Always returning precise element type is incorrect,
1934         // // e.g., element type could be object and array may contain strings
1935         // return TypeKlassPtr::make(TypePtr::Constant, elem, 0);

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1937     // The array's TypeKlassPtr was declared 'precise' or 'not precise'
1938     // according to the element type's subclassing.
1939     return TypeKlassPtr::make(tkls->ptr(), elem, 0/*offset*/);
1940 }
1941 if( klass->is_instance_klass() && tkls->klass_is_exact() &&
1942     (uint)tkls->offset() == Klass::super_offset_in_bytes() + sizeof(oopDesc)
1943     ciKlass* sup = klass->as_instance_klass()->super();
1944     // The field is Klass::_super. Return its (constant) value.
1945     // (Folds up the 2nd indirection in aClassConstant.getSuperClass().)
1946     return sup ? TypeKlassPtr::make(sup) : TypePtr::NULL_PTR;
1947 }
1948 }

1950 // Bailout case
1951 return LoadNode::Value(phase);
1952 }

1954 //-----Identity-----
1955 // To clean up reflective code, simplify k.java_mirror.as_klass to plain k.
1956 // Also feed through the klass in Allocate(...klass...).klass.
1957 Node* LoadKlassNode::Identity( PhaseTransform *phase ) {
1958     return klass_identity_common(phase);
1959 }

1961 Node* LoadNode::klass_identity_common(PhaseTransform *phase ) {
1962     Node* x = LoadNode::Identity(phase);
1963     if (x != this) return x;

1965     // Take apart the address into an oop and an offset.
1966     // Return 'this' if we cannot.
1967     Node* adr = in(MemNode::Address);
1968     intptr_t offset = 0;
1969     Node* base = AddPNode::Ideal_base_and_offset(adr, phase, offset);
1970     if (base == NULL) return this;
1971     const TypeOopPtr* toop = phase->type(adr)->isa_oopptr();
1972     if (toop == NULL) return this;

1974     // We can fetch the klass directly through an AllocateNode.
1975     // This works even if the klass is not constant (clone or newArray).
1976     if (offset == oopDesc::klass_offset_in_bytes()) {
1977         Node* allocated_klass = AllocateNode::Ideal_klass(base, phase);
1978         if (allocated_klass != NULL) {
1979             return allocated_klass;
1980         }
1981     }

1983     // Simplify k.java_mirror.as_klass to plain k, where k is a klassOop.
1984     // Simplify ak.component_mirror.array_klass to plain ak, ak an arrayKlass.
1985     // See inline_native_Class_query for occurrences of these patterns.
1986     // Java Example: x.getClass().isAssignableFrom(y)
1987     // Java Example: Array.newInstance(x.getClass().getComponentType(), n)
1988     //
1989     // This improves reflective code, often making the Class
1990     // mirror go completely dead. (Current exception: Class
1991     // mirrors may appear in debug info, but we could clean them out by
1992     // introducing a new debug info operator for klassOop.java_mirror).
1993     if (toop->isa_instptr() && toop->klass() == phase->C->env()->Class_klass()
1994         && (offset == java_lang_Class::klass_offset_in_bytes() ||
1995             offset == java_lang_Class::array_klass_offset_in_bytes())) {
1996         // We are loading a special hidden field from a Class mirror,
1997         // the field which points to its Klass or arrayKlass metaobject.
1998         if (base->is_Load()) {
1999             Node* adr2 = base->in(MemNode::Address);
2000             const TypeKlassPtr* tkls = phase->type(adr2)->isa_klassptr();
2001             if (tkls != NULL && !tkls->empty()
2002                 && (tkls->klass()->is_instance_klass() ||

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2003     tkls->klass()->is_array_class()
2004     && adr2->is_AddP()
2005     ) {
2006     int mirror_field = Klass::java_mirror_offset_in_bytes();
2007     if (offset == java_lang_Class::array_class_offset_in_bytes()) {
2008         mirror_field = in_bytes(arrayKlass::component_mirror_offset());
2009     }
2010     if (tkls->offset() == mirror_field + (int)sizeof(oopDesc)) {
2011         return adr2->in(AddPNode::Base);
2012     }
2013     }
2014 }
2015 }
2017 return this;
2018 }

2021 //-----Value-----
2022 const Type *LoadNkClassNode::Value( PhaseTransform *phase ) const {
2023     const Type *t = klass_value_common(phase);
2024     if (t == Type::TOP)
2025         return t;
2027     return t->make_narrowoop();
2028 }

2030 //-----Identity-----
2031 // To clean up reflective code, simplify k.java_mirror.as_klass to narrow k.
2032 // Also feed through the klass in Allocate(...klass...)_klass.
2033 Node* LoadNkClassNode::Identity( PhaseTransform *phase ) {
2034     Node *x = klass_identity_common(phase);

2036     const Type *t = phase->type( x );
2037     if( t == Type::TOP ) return x;
2038     if( t->isa_narrowoop() ) return x;

2040     return phase->transform(new ( phase->C, 2) EncodePNode(x, t->make_narrowoop()))
2041 }

2043 //-----Value-----
2044 const Type *LoadRangeNode::Value( PhaseTransform *phase ) const {
2045     // Either input is TOP ==> the result is TOP
2046     const Type *t1 = phase->type( in(MemNode::Memory) );
2047     if( t1 == Type::TOP ) return Type::TOP;
2048     Node *adr = in(MemNode::Address);
2049     const Type *t2 = phase->type( adr );
2050     if( t2 == Type::TOP ) return Type::TOP;
2051     const TypePtr *tp = t2->is_ptr();
2052     if (TypePtr::above_centerline(tp->ptr()) ) return Type::TOP;
2053     const TypeAryPtr *tap = tp->isa_aryptr();
2054     if( !tap ) return_type;
2055     return tap->size();
2056 }

2058 //-----Ideal-----
2059 // Feed through the length in AllocateArray(...length...)_length.
2060 Node *LoadRangeNode::Ideal(PhaseGVN *phase, bool can_reshape) {
2061     Node* p = MemNode::Ideal_common(phase, can_reshape);
2062     if (p) return (p == NodeSentinel) ? NULL : p;

2064     // Take apart the address into an oop and and offset.
2065     // Return 'this' if we cannot.
2066     Node* adr = in(MemNode::Address);
2067     intptr_t offset = 0;
2068     Node* base = AddPNode::Ideal_base_and_offset(adr, phase, offset);

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2069     if (base == NULL) return NULL;
2070     const TypeAryPtr* tary = phase->type(adr)->isa_aryptr();
2071     if (tary == NULL) return NULL;

2073     // We can fetch the length directly through an AllocateArrayNode.
2074     // This works even if the length is not constant (clone or newArray).
2075     if (offset == arrayOopDesc::length_offset_in_bytes()) {
2076         AllocateArrayNode* alloc = AllocateArrayNode::Ideal_array_allocation(base, p
2077         if (alloc != NULL) {
2078             Node* allocated_length = alloc->Ideal_length();
2079             Node* len = alloc->make_ideal_length(tary, phase);
2080             if (allocated_length != len) {
2081                 // New CastII improves on this.
2082                 return len;
2083             }
2084         }
2085     }

2087     return NULL;
2088 }

2090 //-----Identity-----
2091 // Feed through the length in AllocateArray(...length...)_length.
2092 Node* LoadRangeNode::Identity( PhaseTransform *phase ) {
2093     Node* x = LoadINode::Identity(phase);
2094     if (x != this) return x;

2096     // Take apart the address into an oop and and offset.
2097     // Return 'this' if we cannot.
2098     Node* adr = in(MemNode::Address);
2099     intptr_t offset = 0;
2100     Node* base = AddPNode::Ideal_base_and_offset(adr, phase, offset);
2101     if (base == NULL) return this;
2102     const TypeAryPtr* tary = phase->type(adr)->isa_aryptr();
2103     if (tary == NULL) return this;

2105     // We can fetch the length directly through an AllocateArrayNode.
2106     // This works even if the length is not constant (clone or newArray).
2107     if (offset == arrayOopDesc::length_offset_in_bytes()) {
2108         AllocateArrayNode* alloc = AllocateArrayNode::Ideal_array_allocation(base, p
2109         if (alloc != NULL) {
2110             Node* allocated_length = alloc->Ideal_length();
2111             // Do not allow make_ideal_length to allocate a CastII node.
2112             Node* len = alloc->make_ideal_length(tary, phase, false);
2113             if (allocated_length == len) {
2114                 // Return allocated_length only if it would not be improved by a CastII.
2115                 return allocated_length;
2116             }
2117         }
2118     }

2120     return this;
2122 }

2124 //-----StoreNode:make-----
2125 // Polymorphic factory method:
2126 StoreNode* StoreNode::make( PhaseGVN& gvn, Node* ctl, Node* mem, Node* adr, cons
2127 Compile* C = gvn.C;
2128 assert( C->get_alias_index(adr_type) != Compile::AliasIdxRaw ||
2129         ctl != NULL, "raw memory operations should have control edge");

2132     switch (bt) {
2133     case T_BOOLEAN:
2134     case T_BYTE: return new (C, 4) StoreBNode(ctl, mem, adr, adr_type, val);

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2135 case T_INT: return new (C, 4) StoreINode(ctl, mem, adr, adr_type, val);
2136 case T_CHAR:
2137 case T_SHORT: return new (C, 4) StoreCNode(ctl, mem, adr, adr_type, val);
2138 case T_LONG: return new (C, 4) StoreLNode(ctl, mem, adr, adr_type, val);
2139 case T_FLOAT: return new (C, 4) StoreFNode(ctl, mem, adr, adr_type, val);
2140 case T_DOUBLE: return new (C, 4) StoreDNode(ctl, mem, adr, adr_type, val);
2141 case T_ADDRESS:
2142 case T_OBJECT:
2143 #ifdef LP64
2144 if (adr->bottom_type()->is_ptr_to_narrowoop() ||
2145 (UseCompressedOops && val->bottom_type()->isa_klassptr() &&
2146 adr->bottom_type()->isa_rawptr())) {
2147 val = gvn.transform(new (C, 2) EncodePNode(val, val->bottom_type()->make_n
2148 return new (C, 4) StoreNNode(ctl, mem, adr, adr_type, val);
2149 } else
2150 #endif
2151 {
2152 return new (C, 4) StorePNode(ctl, mem, adr, adr_type, val);
2153 }
2154 }
2155 ShouldNotReachHere();
2156 return (StoreNode*)NULL;
2157 }

2159 StoreLNode* StoreLNode::make_atomic(Compile *C, Node* ctl, Node* mem, Node* adr,
2160 bool require_atomic = true;
2161 return new (C, 4) StoreLNode(ctl, mem, adr, adr_type, val, require_atomic);
2162 }

2165 //-----bottom_type-----
2166 const Type *StoreNode::bottom_type() const {
2167 return Type::MEMORY;
2168 }

2170 //-----hash-----
2171 uint StoreNode::hash() const {
2172 // unroll addition of interesting fields
2173 //return (uintptr_t)in(Control) + (uintptr_t)in(Memory) + (uintptr_t)in(Addres

2175 // Since they are not commoned, do not hash them:
2176 return NO_HASH;
2177 }

2179 //-----Ideal-----
2180 // Change back-to-back Store(, p, x) -> Store(m, p, y) to Store(m, p, x).
2181 // When a store immediately follows a relevant allocation/initialization,
2182 // try to capture it into the initialization, or hoist it above.
2183 Node *StoreNode::Ideal(PhaseGVN *phase, bool can_reshape) {
2184 Node* p = MemNode::Ideal_common(phase, can_reshape);
2185 if (p) return (p == NodeSentinel) ? NULL : p;

2187 Node* mem = in(MemNode::Memory);
2188 Node* address = in(MemNode::Address);

2190 // Back-to-back stores to same address? Fold em up. Generally
2191 // unsafe if I have intervening uses... Also disallowed for StoreCM
2192 // since they must follow each StoreP operation. Redundant StoreCMs
2193 // are eliminated just before matching in final_graph_reshape.
2194 if (mem->is_Store() && phase->eqv_unCast(mem->in(MemNode::Address), address) &
2195 mem->Opcode() != Op_StoreCM) {
2196 // Looking at a dead closed cycle of memory?
2197 assert(mem != mem->in(MemNode::Memory), "dead loop in StoreNode::Ideal");

2199 assert(Opcode() == mem->Opcode() ||
2200 phase->C->get_alias_index(adr_type()) == Compile::AliasIdxRaw,

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2201 "no mismatched stores, except on raw memory");

2203 if (mem->outcnt() == 1 && // check for intervening uses
2204 mem->as_Store()->memory_size() <= this->memory_size()) {
2205 // If anybody other than 'this' uses 'mem', we cannot fold 'mem' away.
2206 // For example, 'mem' might be the final state at a conditional return.
2207 // Or, 'mem' might be used by some node which is live at the same time
2208 // 'this' is live, which might be unschedulable. So, require exactly
2209 // ONE user, the 'this' store, until such time as we clone 'mem' for
2210 // each of 'mem's uses (thus making the exactly-1-user-rule hold true).
2211 if (can_reshape) { // (%% is this an anachronism?)
2212 set_req_X(MemNode::Memory, mem->in(MemNode::Memory),
2213 phase->is_IterGVN());
2214 } else {
2215 // It's OK to do this in the parser, since DU info is always accurate,
2216 // and the parser always refers to nodes via SafePointNode maps.
2217 set_req(MemNode::Memory, mem->in(MemNode::Memory));
2218 }
2219 return this;
2220 }
2221 }

2223 // Capture an unaliased, unconditional, simple store into an initializer.
2224 // Or, if it is independent of the allocation, hoist it above the allocation.
2225 if (ReduceFieldZeroing && /*can_reshape &&*/
2226 mem->is_Proj() && mem->in(0)->is_Initialize()) {
2227 InitializeNode* init = mem->in(0)->as_Initialize();
2228 intptr_t offset = init->can_capture_store(this, phase);
2229 if (offset > 0) {
2230 Node* moved = init->capture_store(this, offset, phase);
2231 // If the InitializeNode captured me, it made a raw copy of me,
2232 // and I need to disappear.
2233 if (moved != NULL) {
2234 // %% hack to ensure that Ideal returns a new node:
2235 mem = MergeMemNode::make(phase->C, mem);
2236 return mem; // fold me away
2237 }
2238 }
2239 }

2241 return NULL; // No further progress
2242 }

2244 //-----Value-----
2245 const Type *StoreNode::Value(PhaseTransform *phase) const {
2246 // Either input is TOP => the result is TOP
2247 const Type *t1 = phase->type(in(MemNode::Memory));
2248 if (t1 == Type::TOP) return Type::TOP;
2249 const Type *t2 = phase->type(in(MemNode::Address));
2250 if (t2 == Type::TOP) return Type::TOP;
2251 const Type *t3 = phase->type(in(MemNode::ValueIn));
2252 if (t3 == Type::TOP) return Type::TOP;
2253 return Type::MEMORY;
2254 }

2256 //-----Identity-----
2257 // Remove redundant stores:
2258 // Store(m, p, Load(m, p)) changes to m.
2259 // Store(m, p, x) -> Store(m, p, x) changes to Store(m, p, x).
2260 Node *StoreNode::Identity(PhaseTransform *phase) {
2261 Node* mem = in(MemNode::Memory);
2262 Node* adr = in(MemNode::Address);
2263 Node* val = in(MemNode::ValueIn);

2265 // Load then Store? Then the Store is useless
2266 if (val->is_Load() &&

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2267     phase->eqv_uncanst( val->in(MemNode::Address), adr ) &&
2268     phase->eqv_uncanst( val->in(MemNode::Memory), mem ) &&
2269     val->as_Load()->store_Opcode() == Opcode() ) {
2270     return mem;
2271 }

2273 // Two stores in a row of the same value?
2274 if (mem->is_Store() &&
2275     phase->eqv_uncanst( mem->in(MemNode::Address), adr ) &&
2276     phase->eqv_uncanst( mem->in(MemNode::ValueIn), val ) &&
2277     mem->Opcode() == Opcode() ) {
2278     return mem;
2279 }

2281 // Store of zero anywhere into a freshly-allocated object?
2282 // Then the store is useless.
2283 // (It must already have been captured by the InitializeNode.)
2284 if (ReduceFieldZeroing && phase->type(val)->is_zero_type()) {
2285     // a newly allocated object is already all-zeroes everywhere
2286     if (mem->is_Proj() && mem->in(0)->is_Allocate()) {
2287         return mem;
2288     }

2290     // the store may also apply to zero-bits in an earlier object
2291     Node* prev_mem = find_previous_store(phase);
2292     // Steps (a), (b): Walk past independent stores to find an exact match.
2293     if (prev_mem != NULL) {
2294         Node* prev_val = can_see_stored_value(prev_mem, phase);
2295         if (prev_val != NULL && phase->eqv(prev_val, val)) {
2296             // prev_val and val might differ by a cast; it would be good
2297             // to keep the more informative of the two.
2298             return mem;
2299         }
2300     }
2301 }

2303 return this;
2304 }

2306 //-----match_edge-----
2307 // Do we Match on this edge index or not? Match only memory & value
2308 uint StoreNode::match_edge(uint idx) const {
2309     return idx == MemNode::Address || idx == MemNode::ValueIn;
2310 }

2312 //-----cmp-----
2313 // Do not common stores up together. They generally have to be split
2314 // back up anyways, so do not bother.
2315 uint StoreNode::cmp( const Node &n ) const {
2316     return (&n == this); // Always fail except on self
2317 }

2319 //-----Ideal_masked_input-----
2320 // Check for a useless mask before a partial-word store
2321 // (StoreB ... (AndI valIn conIa) )
2322 // If (conIa & mask == mask) this simplifies to
2323 // (StoreB ... (valIn) )
2324 Node *StoreNode::Ideal_masked_input(PhaseGVN *phase, uint mask) {
2325     Node *val = in(MemNode::ValueIn);
2326     if (val->Opcode() == Op_AndI ) {
2327         const TypeInt *t = phase->type( val->in(2) )->isa_int();
2328         if( t && t->is_con() && (t->get_con() & mask) == mask ) {
2329             set_req(MemNode::ValueIn, val->in(1));
2330             return this;
2331         }
2332     }

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2333     return NULL;
2334 }

2337 //-----Ideal_sign_extended_input-----
2338 // Check for useless sign-extension before a partial-word store
2339 // (StoreB ... (RShiftI _ (LShiftI _ valIn conIL) conIR) )
2340 // If (conIL == conIR && conIR <= num_bits) this simplifies to
2341 // (StoreB ... (valIn) )
2342 Node *StoreNode::Ideal_sign_extended_input(PhaseGVN *phase, int num_bits) {
2343     Node *val = in(MemNode::ValueIn);
2344     if( val->Opcode() == Op_RShiftI ) {
2345         const TypeInt *t = phase->type( val->in(2) )->isa_int();
2346         if( t && t->is_con() && (t->get_con() <= num_bits) ) {
2347             Node *shl = val->in(1);
2348             if( shl->Opcode() == Op_LShiftI ) {
2349                 const TypeInt *t2 = phase->type( shl->in(2) )->isa_int();
2350                 if( t2 && t2->is_con() && (t2->get_con() == t->get_con()) ) {
2351                     set_req(MemNode::ValueIn, shl->in(1));
2352                     return this;
2353                 }
2354             }
2355         }
2356     }
2357     return NULL;
2358 }

2360 //-----value_never_loaded-----
2361 // Determine whether there are any possible loads of the value stored.
2362 // For simplicity, we actually check if there are any loads from the
2363 // address stored to, not just for loads of the value stored by this node.
2364 //
2365 bool StoreNode::value_never_loaded( PhaseTransform *phase) const {
2366     Node *adr = in(Address);
2367     const TypeOopPtr *adr_oop = phase->type(adr)->isa_oopptr();
2368     if (adr_oop == NULL)
2369         return false;
2370     if (!adr_oop->is_known_instance_field())
2371         return false; // if not a distinct instance, there may be aliases of the add
2372     for (DUIterator_Fast imax, i = adr->fast_outs(imax); i < imax; i++) {
2373         Node *use = adr->fast_out(i);
2374         int opc = use->Opcode();
2375         if (use->is_Load() || use->is_LoadStore()) {
2376             return false;
2377         }
2378     }
2379     return true;
2380 }

2382 //=====
2383 //-----Ideal-----
2384 // If the store is from an AND mask that leaves the low bits untouched, then
2385 // we can skip the AND operation. If the store is from a sign-extension
2386 // (a left shift, then right shift) we can skip both.
2387 Node *StoreNode::Ideal(PhaseGVN *phase, bool can_reshape){
2388     Node *progress = StoreNode::Ideal_masked_input(phase, 0xFF);
2389     if( progress != NULL ) return progress;

2391     progress = StoreNode::Ideal_sign_extended_input(phase, 24);
2392     if( progress != NULL ) return progress;

2394     // Finally check the default case
2395     return StoreNode::Ideal(phase, can_reshape);
2396 }

2398 //=====

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```

2399 //-----Ideal-----
2400 // If the store is from an AND mask that leaves the low bits untouched, then
2401 // we can skip the AND operation
2402 Node *StoreCNode::Ideal(PhaseGVN *phase, bool can_reshape){
2403     Node *progress = StoreNode::Ideal_masked_input(phase, 0xFFFF);
2404     if( progress != NULL ) return progress;

2406     progress = StoreNode::Ideal_sign_extended_input(phase, 16);
2407     if( progress != NULL ) return progress;

2409     // Finally check the default case
2410     return StoreNode::Ideal(phase, can_reshape);
2411 }

2413 //=====
2414 //-----Identity-----
2415 Node *StoreCMNode::Identity( PhaseTransform *phase ) {
2416     // No need to card mark when storing a null ptr
2417     Node* my_store = in(MemNode::OopStore);
2418     if (my_store->is_Store()) {
2419         const Type *t1 = phase->type( my_store->in(MemNode::ValueIn) );
2420         if( t1 == TypePtr::NULL_PTR ) {
2421             return in(MemNode::Memory);
2422         }
2423     }
2424     return this;
2425 }

2427 //=====
2428 //-----Ideal-----
2429 Node *StoreCMNode::Ideal(PhaseGVN *phase, bool can_reshape){
2430     Node* progress = StoreNode::Ideal(phase, can_reshape);
2431     if( progress != NULL ) return progress;

2433     Node* my_store = in(MemNode::OopStore);
2434     if (my_store->is_MergeMem()) {
2435         Node* mem = my_store->as_MergeMem()->memory_at(oop_alias_idx());
2436         set_req(MemNode::OopStore, mem);
2437         return this;
2438     }

2440     return NULL;
2441 }

2443 //-----Value-----
2444 const Type *StoreCMNode::Value( PhaseTransform *phase ) const {
2445     // Either input is TOP ==> the result is TOP
2446     const Type *t = phase->type( in(MemNode::Memory) );
2447     if( t == Type::TOP ) return Type::TOP;
2448     t = phase->type( in(MemNode::Address) );
2449     if( t == Type::TOP ) return Type::TOP;
2450     t = phase->type( in(MemNode::ValueIn) );
2451     if( t == Type::TOP ) return Type::TOP;
2452     // If extra input is TOP ==> the result is TOP
2453     t = phase->type( in(MemNode::OopStore) );
2454     if( t == Type::TOP ) return Type::TOP;

2456     return StoreNode::Value( phase );
2457 }

2460 //=====
2461 //-----SCMemProjNode-----
2462 const Type * SCMemProjNode::Value( PhaseTransform *phase ) const
2463 {
2464     return bottom_type();

```

```

2465 }

2467 //=====
2468 LoadStoreNode::LoadStoreNode( Node *c, Node *mem, Node *adr, Node *val, Node *ex
2469     init_req(MemNode::Control, c );
2470     init_req(MemNode::Memory, mem);
2471     init_req(MemNode::Address, adr);
2472     init_req(MemNode::ValueIn, val);
2473     init_req( ExpectedIn, ex );
2474     init_class_id(Class_LoadStore);

2476 }

2478 //=====
2479 //-----adr_type-----
2480 // Do we Match on this edge index or not? Do not match memory
2481 const TypePtr* ClearArrayNode::adr_type() const {
2482     Node *adr = in(3);
2483     return MemNode::calculate_adr_type(adr->bottom_type());
2484 }

2486 //-----match_edge-----
2487 // Do we Match on this edge index or not? Do not match memory
2488 uint ClearArrayNode::match_edge(uint idx) const {
2489     return idx > 1;
2490 }

2492 //-----Identity-----
2493 // Clearing a zero length array does nothing
2494 Node *ClearArrayNode::Identity( PhaseTransform *phase ) {
2495     return phase->type(in(2))->higher_equal(TypeX::ZERO) ? in(1) : this;
2496 }

2498 //-----Idealize-----
2499 // Clearing a short array is faster with stores
2500 Node *ClearArrayNode::Ideal(PhaseGVN *phase, bool can_reshape){
2501     const int unit = BytesPerLong;
2502     const TypeX* t = phase->type(in(2))->isa_intptr_t();
2503     if (!t) return NULL;
2504     if (!t->is_con()) return NULL;
2505     intptr_t raw_count = t->get_con();
2506     intptr_t size = raw_count;
2507     if (!Matcher::init_array_count_is_in_bytes) size *= unit;
2508     // Clearing nothing uses the Identity call.
2509     // Negative clears are possible on dead ClearArrays
2510     // (see jck test stmt114.stmt11402.val).
2511     if (size <= 0 || size % unit != 0) return NULL;
2512     intptr_t count = size / unit;
2513     // Length too long; use fast hardware clear
2514     if (size > Matcher::init_array_short_size) return NULL;
2515     Node *mem = in(1);
2516     if( phase->type(mem)==Type::TOP ) return NULL;
2517     Node *adr = in(3);
2518     const Type* at = phase->type(adr);
2519     if( at==Type::TOP ) return NULL;
2520     const TypePtr* atp = at->isa_ptr();
2521     // adjust atp to be the correct array element address type
2522     if (atp == NULL) atp = TypePtr::BOTTOM;
2523     else atp = atp->add_offset(Type::OffsetBot);
2524     // Get base for derived pointer purposes
2525     if( adr->Opcode() != Op_AddP ) Unimplemented();
2526     Node *base = adr->in(1);

2528     Node *zero = phase->makecon(TypeLong::ZERO);
2529     Node *off = phase->MakeConX(BytesPerLong);
2530     mem = new (phase->C, 4) StoreLNode(in(0),mem,adr,atp,zero);

```



```

2531 count--;
2532 while( count-- ) {
2533     mem = phase->transform(mem);
2534     adr = phase->transform(new (phase->C, 4) AddPNode(base,adr,off));
2535     mem = new (phase->C, 4) StoreLNode(in(0),mem,adr,atp,zero);
2536 }
2537 return mem;
2538 }

2540 //-----step_through-----
2541 // Return allocation input memory edge if it is different instance
2542 // or itself if it is the one we are looking for.
2543 bool ClearArrayNode::step_through(Node** np, uint instance_id, PhaseTransform* p
2544 Node* n = *np;
2545 assert(n->is_ClearArray(), "sanity");
2546 intptr_t offset;
2547 AllocateNode* alloc = AllocateNode::Ideal_allocation(n->in(3), phase, offset);
2548 // This method is called only before Allocate nodes are expanded during
2549 // macro nodes expansion. Before that ClearArray nodes are only generated
2550 // in LibraryCallKit::generate_arraycopy() which follows allocations.
2551 assert(alloc != NULL, "should have allocation");
2552 if (alloc->idx == instance_id) {
2553     // Can not bypass initialization of the instance we are looking for.
2554     return false;
2555 }
2556 // Otherwise skip it.
2557 InitializeNode* init = alloc->initialization();
2558 if (init != NULL)
2559     *np = init->in(TypeFunc::Memory);
2560 else
2561     *np = alloc->in(TypeFunc::Memory);
2562 return true;
2563 }

2565 //-----clear_memory-----
2566 // Generate code to initialize object storage to zero.
2567 Node* ClearArrayNode::clear_memory(Node* ctl, Node* mem, Node* dest,
2568     intptr_t start_offset,
2569     Node* end_offset,
2570     PhaseGVN* phase) {
2571     Compile* C = phase->C;
2572     intptr_t offset = start_offset;

2574     int unit = BytesPerLong;
2575     if ((offset % unit) != 0) {
2576         Node* adr = new (C, 4) AddPNode(dest, dest, phase->MakeConX(offset));
2577         adr = phase->transform(adr);
2578         const TypePtr* atp = TypeRawPtr::BOTTOM;
2579         mem = StoreNode::make(*phase, ctl, mem, adr, atp, phase->zerocon(T_INT), T_I
2580         mem = phase->transform(mem);
2581         offset += BytesPerInt;
2582     }
2583     assert((offset % unit) == 0, "");

2585     // Initialize the remaining stuff, if any, with a ClearArray.
2586     return clear_memory(ctl, mem, dest, phase->MakeConX(offset), end_offset, phase
2587 }

2589 Node* ClearArrayNode::clear_memory(Node* ctl, Node* mem, Node* dest,
2590     Node* start_offset,
2591     Node* end_offset,
2592     PhaseGVN* phase) {
2593     if (start_offset == end_offset) {
2594         // nothing to do
2595         return mem;
2596     }

```

```

2598     Compile* C = phase->C;
2599     int unit = BytesPerLong;
2600     Node* zbase = start_offset;
2601     Node* zend = end_offset;

2603     // Scale to the unit required by the CPU:
2604     if (!Matcher::init_array_count_is_in_bytes) {
2605         Node* shift = phase->intcon(exact_log2(unit));
2606         zbase = phase->transform( new(C,3) URShiftXNode(zbase, shift) );
2607         zend = phase->transform( new(C,3) URShiftXNode(zend, shift) );
2608     }

2610     Node* zsize = phase->transform( new(C,3) SubXNode(zend, zbase) );
2611     Node* zinit = phase->zerocon((unit == BytesPerLong) ? T_LONG : T_INT);

2613     // Bulk clear double-words
2614     Node* adr = phase->transform( new(C,4) AddPNode(dest, dest, start_offset) );
2615     mem = new (C, 4) ClearArrayNode(ctl, mem, zsize, adr);
2616     return phase->transform(mem);
2617 }

2619 Node* ClearArrayNode::clear_memory(Node* ctl, Node* mem, Node* dest,
2620     intptr_t start_offset,
2621     intptr_t end_offset,
2622     PhaseGVN* phase) {
2623     if (start_offset == end_offset) {
2624         // nothing to do
2625         return mem;
2626     }

2628     Compile* C = phase->C;
2629     assert((end_offset % BytesPerInt) == 0, "odd end offset");
2630     intptr_t done_offset = end_offset;
2631     if ((done_offset % BytesPerLong) != 0) {
2632         done_offset -= BytesPerInt;
2633     }
2634     if (done_offset > start_offset) {
2635         mem = clear_memory(ctl, mem, dest,
2636             start_offset, phase->MakeConX(done_offset), phase);
2637     }
2638     if (done_offset < end_offset) { // emit the final 32-bit store
2639         Node* adr = new (C, 4) AddPNode(dest, dest, phase->MakeConX(done_offset));
2640         adr = phase->transform(adr);
2641         const TypePtr* atp = TypeRawPtr::BOTTOM;
2642         mem = StoreNode::make(*phase, ctl, mem, adr, atp, phase->zerocon(T_INT), T_I
2643         mem = phase->transform(mem);
2644         done_offset += BytesPerInt;
2645     }
2646     assert(done_offset == end_offset, "");
2647     return mem;
2648 }

2650 //=====
2651 // Do not match memory edge.
2652 uint StrIntrinsicNode::match_edge(uint idx) const {
2653     return idx == 2 || idx == 3;
2654 }

2656 //-----Ideal-----
2657 // Return a node which is more "ideal" than the current node. Strip out
2658 // control copies
2659 Node* StrIntrinsicNode::Ideal(PhaseGVN* phase, bool can_reshape) {
2660     if (remove_dead_region(phase, can_reshape)) return this;

2662     if (can_reshape) {

```

```

2663 Node* mem = phase->transform(in(MemNode::Memory));
2664 // If transformed to a MergeMem, get the desired slice
2665 uint alias_idx = phase->C->get_alias_index(adr_type());
2666 mem = mem->is_MergeMem() ? mem->as_MergeMem()->memory_at(alias_idx) : mem;
2667 if (mem != in(MemNode::Memory)) {
2668     set_req(MemNode::Memory, mem);
2669     return this;
2670 }
2671 }
2672 return NULL;
2673 }

2675 //=====
2676 MemBarNode::MemBarNode(Compile* C, int alias_idx, Node* precedent)
2677 : MultiNode(TypeFunc::Parms + (precedent == NULL? 0: 1)),
2678   _adr_type(C->get_adr_type(alias_idx))
2679 {
2680     init_class_id(Class_MemBar);
2681     Node* top = C->top();
2682     init_req(TypeFunc::I_O, top);
2683     init_req(TypeFunc::FramePtr, top);
2684     init_req(TypeFunc::ReturnAdr, top);
2685     if (precedent != NULL)
2686         init_req(TypeFunc::Parms, precedent);
2687 }

2689 //-----cmp-----
2690 uint MemBarNode::hash() const { return NO_HASH; }
2691 uint MemBarNode::cmp( const Node &n ) const {
2692     return (&n == this); // Always fail except on self
2693 }

2695 //-----make-----
2696 MemBarNode* MemBarNode::make(Compile* C, int opcode, int atp, Node* pn) {
2697     int len = Precedent + (pn == NULL? 0: 1);
2698     switch (opcode) {
2699     case Op_MemBarAcquire: return new(C, len) MemBarAcquireNode(C, atp, pn);
2700     case Op_MemBarRelease: return new(C, len) MemBarReleaseNode(C, atp, pn);
2701     case Op_MemBarAcquireLock: return new(C, len) MemBarAcquireLockNode(C, atp, p
2702     case Op_MemBarReleaseLock: return new(C, len) MemBarReleaseLockNode(C, atp, p
2703     case Op_MemBarVolatile: return new(C, len) MemBarVolatileNode(C, atp, pn);
2704     case Op_MemBarCPUOrder: return new(C, len) MemBarCPUOrderNode(C, atp, pn);
2705     case Op_Initialize: return new(C, len) InitializeNode(C, atp, pn);
2706     default: ShouldNotReachHere(); return NULL;
2707     }
2708 }

2710 //-----Ideal-----
2711 // Return a node which is more "ideal" than the current node. Strip out
2712 // control copies
2713 Node* MemBarNode::Ideal(PhaseGVN* phase, bool can_reshape) {
2714     if (remove_dead_region(phase, can_reshape)) replaced this;

2716     // Eliminate volatile MemBars for scalar replaced objects.
2717     if (can_reshape && req() == (Precedent+1) &&
2718         (Opcode() == Op_MemBarAcquire || Opcode() == Op_MemBarVolatile)) {
2719         // Volatile field loads and stores.
2720         Node* my_mem = in(MemBarNode::Precedent);
2721         if (my_mem != NULL && my_mem->is_Mem()) {
2722             const TypeOopPtr* t_oop = my_mem->in(MemNode::Address)->bottom_type()->isa
2723             // Check for scalar replaced object reference.
2724             if( t_oop != NULL && t_oop->is_known_instance_field() &&
2725                 t_oop->offset() != Type::OffsetBot &&
2726                 t_oop->offset() != Type::OffsetTop) {
2727                 // Replace MemBar projections by its inputs.
2728                 PhaseIterGVN* igvn = phase->is_IterGVN();

```

```

2729     igvn->replace_node(proj_out(TypeFunc::Memory), in(TypeFunc::Memory));
2730     igvn->replace_node(proj_out(TypeFunc::Control), in(TypeFunc::Control));
2731     // Must return either the original node (now dead) or a new node
2732     // (Do not return a top here, since that would break the uniqueness of t
2733     return new (phase->C, 1) ConINode(TypeInt::ZERO);
2734     }
2735     }
2736     }
2737     return NULL;
2738 }

2740 //-----Value-----
2741 const Type* MemBarNode::Value( PhaseTransform* phase ) const {
2742     if( !in(0) ) return Type::TOP;
2743     if( phase->type(in(0)) == Type::TOP )
2744         return Type::TOP;
2745     return TypeTuple::MEMBAR;
2746 }

2748 //-----match-----
2749 // Construct projections for memory.
2750 Node* MemBarNode::match( const ProjNode* proj, const Matcher* m ) {
2751     switch (proj->_con) {
2752     case TypeFunc::Control:
2753     case TypeFunc::Memory:
2754         return new (m->C, 1) MachProjNode(this, proj->_con, RegMask::Empty, MachProjNod
2755     }
2756     ShouldNotReachHere();
2757     return NULL;
2758 }

2760 //=====InitializeNode=====
2761 // SUMMARY:
2762 // This node acts as a memory barrier on raw memory, after some raw stores.
2763 // The 'cooked' oop value feeds from the Initialize, not the Allocation.
2764 // The Initialize can 'capture' suitably constrained stores as raw inits.
2765 // It can coalesce related raw stores into larger units (called 'tiles').
2766 // It can avoid zeroing new storage for memory units which have raw inits.
2767 // At macro-expansion, it is marked 'complete', and does not optimize further.
2768 //
2769 // EXAMPLE:
2770 // The object 'new short[2]' occupies 16 bytes in a 32-bit machine.
2771 //   ctl = incoming control; mem* = incoming memory
2772 // (Note: A star * on a memory edge denotes I/O and other standard edges.)
2773 // First allocate uninitialized memory and fill in the header:
2774 //   alloc = (Allocate ctl mem* 16 #short[.k]klass ...)
2775 //   ctl := alloc.Control; mem* := alloc.Memory*
2776 //   rawmem = alloc.Memory; rawoop = alloc.RawAddress
2777 // Then initialize to zero the non-header parts of the raw memory block:
2778 //   init = (Initialize alloc.Control alloc.Memory* alloc.RawAddress)
2779 //   ctl := init.Control; mem.SLICE(#short[*]) := init.Memory
2780 // After the initialize node executes, the object is ready for service:
2781 //   oop := (CheckCastPP init.Control alloc.RawAddress #short[.k])
2782 // Suppose its body is immediately initialized as {1,2}:
2783 //   store1 = (StoreC init.Control init.Memory (+ oop 12) 1)
2784 //   store2 = (StoreC init.Control store1 (+ oop 14) 2)
2785 //   mem.SLICE(#short[*]) := store2
2786 //
2787 // DETAILS:
2788 // An InitializeNode collects and isolates object initialization after
2789 // an AllocateNode and before the next possible safepoint. As a
2790 // memory barrier (MemBarNode), it keeps critical stores from drifting
2791 // down past any safepoint or any publication of the allocation.
2792 // Before this barrier, a newly-allocated object may have uninitialized bits.
2793 // After this barrier, it may be treated as a real oop, and GC is allowed.
2794 //

```

```

2795 // The semantics of the InitializeNode include an implicit zeroing of
2796 // the new object from object header to the end of the object.
2797 // (The object header and end are determined by the AllocateNode.)
2798 //
2799 // Certain stores may be added as direct inputs to the InitializeNode.
2800 // These stores must update raw memory, and they must be to addresses
2801 // derived from the raw address produced by AllocateNode, and with
2802 // a constant offset. They must be ordered by increasing offset.
2803 // The first one is at in(RawStores), the last at in(req()-1).
2804 // Unlike most memory operations, they are not linked in a chain,
2805 // but are displayed in parallel as users of the rawmem output of
2806 // the allocation.
2807 //
2808 // (See comments in InitializeNode::capture_store, which continue
2809 // the example given above.)
2810 //
2811 // When the associated Allocate is macro-expanded, the InitializeNode
2812 // may be rewritten to optimize collected stores. A ClearArrayNode
2813 // may also be created at that point to represent any required zeroing.
2814 // The InitializeNode is then marked 'complete', prohibiting further
2815 // capturing of nearby memory operations.
2816 //
2817 // During macro-expansion, all captured initializations which store
2818 // constant values of 32 bits or smaller are coalesced (if advantageous)
2819 // into larger 'tiles' 32 or 64 bits. This allows an object to be
2820 // initialized in fewer memory operations. Memory words which are
2821 // covered by neither tiles nor non-constant stores are pre-zeroed
2822 // by explicit stores of zero. (The code shape happens to do all
2823 // zeroing first, then all other stores, with both sequences occurring
2824 // in order of ascending offsets.)
2825 //
2826 // Alternatively, code may be inserted between an AllocateNode and its
2827 // InitializeNode, to perform arbitrary initialization of the new object.
2828 // E.g., the object copying intrinsics insert complex data transfers here.
2829 // The initialization must then be marked as 'complete' disable the
2830 // built-in zeroing semantics and the collection of initializing stores.
2831 //
2832 // While an InitializeNode is incomplete, reads from the memory state
2833 // produced by it are optimizable if they match the control edge and
2834 // new oop address associated with the allocation/initialization.
2835 // They return a stored value (if the offset matches) or else zero.
2836 // A write to the memory state, if it matches control and address,
2837 // and if it is to a constant offset, may be 'captured' by the
2838 // InitializeNode. It is cloned as a raw memory operation and rewired
2839 // inside the initialization, to the raw oop produced by the allocation.
2840 // Operations on addresses which are provably distinct (e.g., to
2841 // other AllocateNodes) are allowed to bypass the initialization.
2842 //
2843 // The effect of all this is to consolidate object initialization
2844 // (both arrays and non-arrays, both piecewise and bulk) into a
2845 // single location, where it can be optimized as a unit.
2846 //
2847 // Only stores with an offset less than TrackedInitializationLimit words
2848 // will be considered for capture by an InitializeNode. This puts a
2849 // reasonable limit on the complexity of optimized initializations.

```

```

2851 //-----InitializeNode-----
2852 InitializeNode::InitializeNode(Compile* C, int adr_type, Node* rawoop)
2853   : _is_complete(false),
2854     MemBarNode(C, adr_type, rawoop)
2855 {
2856   init_class_id(Class_Initialize);

```

```

2858   assert(adr_type == Compile::AliasIdxRaw, "only valid atp");
2859   assert(in(RawAddress) == rawoop, "proper init");
2860   // Note: allocation() can be NULL, for secondary initialization barriers

```

```

2861 }

```

```

2863 // Since this node is not matched, it will be processed by the
2864 // register allocator. Declare that there are no constraints
2865 // on the allocation of the RawAddress edge.
2866 const RegMask &InitializeNode::in_RegMask(uint idx) const {
2867   // This edge should be set to top, by the set_complete. But be conservative.
2868   if (idx == InitializeNode::RawAddress)
2869     return *(Compile::current()->matcher()->idealreg2spillmask[in(idx)->ideal_re
2870   return RegMask::Empty;
2871 }

```

```

2873 Node* InitializeNode::memory(uint alias_idx) {
2874   Node* mem = in(Memory);
2875   if (mem->is_MergeMem()) {
2876     return mem->as_MergeMem()->memory_at(alias_idx);
2877   } else {
2878     // incoming raw memory is not split
2879     return mem;
2880   }
2881 }

```

```

2883 bool InitializeNode::is_non_zero() {
2884   if (is_complete()) return false;
2885   remove_extra_zeroes();
2886   return (req() > RawStores);
2887 }

```

```

2889 void InitializeNode::set_complete(PhaseGVN* phase) {
2890   assert(!is_complete(), "caller responsibility");
2891   _is_complete = true;

```

```

2893   // After this node is complete, it contains a bunch of
2894   // raw-memory initializations. There is no need for
2895   // it to have anything to do with non-raw memory effects.
2896   // Therefore, tell all non-raw users to re-optimize themselves,
2897   // after skipping the memory effects of this initialization.
2898   PhaseIterGVN* igvn = phase->is_IterGVN();
2899   if (igvn) igvn->add_users_to_worklist(this);
2900 }

```

```

2902 // convenience function
2903 // return false if the init contains any stores already
2904 bool AllocateNode::maybe_set_complete(PhaseGVN* phase) {
2905   InitializeNode* init = initialization();
2906   if (init == NULL || init->is_complete()) return false;
2907   init->remove_extra_zeroes();
2908   // for now, if this allocation has already collected any inits, bail:
2909   if (init->is_non_zero()) return false;
2910   init->set_complete(phase);
2911   return true;
2912 }

```

```

2914 void InitializeNode::remove_extra_zeroes() {
2915   if (req() == RawStores) return;
2916   Node* zmem = zero_memory();
2917   uint fill = RawStores;
2918   for (uint i = fill; i < req(); i++) {
2919     Node* n = in(i);
2920     if (n->is_top() || n == zmem) continue; // skip
2921     if (fill < i) set_req(fill, n); // compact
2922     ++fill;
2923   }
2924   // delete any empty spaces created:
2925   while (fill < req()) {
2926     del_req(fill);

```

```

2927 }
2928 }

2930 // Helper for remembering which stores go with which offsets.
2931 intptr_t InitializeNode::get_store_offset(Node* st, PhaseTransform* phase) {
2932     if (!st->is_Store()) return -1; // can happen to dead code via subsume_node
2933     intptr_t offset = -1;
2934     Node* base = AddPNode::Ideal_base_and_offset(st->in(MemNode::Address),
2935         phase, offset);
2936     if (base == NULL) return -1; // something is dead,
2937     if (offset < 0) return -1; // dead, dead
2938     return offset;
2939 }

2941 // Helper for proving that an initialization expression is
2942 // "simple enough" to be folded into an object initialization.
2943 // Attempts to prove that a store's initial value 'n' can be captured
2944 // within the initialization without creating a vicious cycle, such as:
2945 // { Foo p = new Foo(); p.next = p; }
2946 // True for constants and parameters and small combinations thereof.
2947 bool InitializeNode::detect_init_independence(Node* n,
2948     bool st_is_pinned,
2949     int& count) {
2950     if (n == NULL) return true; // (can this really happen?)
2951     if (n->is_Proj()) n = n->in(0);
2952     if (n == this) return false; // found a cycle
2953     if (n->is_Con()) return true;
2954     if (n->is_Start()) return true; // params, etc., are OK
2955     if (n->is_Root()) return true; // even better

2957     Node* ctl = n->in(0);
2958     if (ctl != NULL && !ctl->is_top()) {
2959         if (ctl->is_Proj()) ctl = ctl->in(0);
2960         if (ctl == this) return false;

2962         // If we already know that the enclosing memory op is pinned right after
2963         // the init, then any control flow that the store has picked up
2964         // must have preceded the init, or else be equal to the init.
2965         // Even after loop optimizations (which might change control edges)
2966         // a store is never pinned *before* the availability of its inputs.
2967         if (!MemNode::all_controls_dominated(n, this))
2968             return false; // failed to prove a good control

2970     }

2972     // Check data edges for possible dependencies on 'this'.
2973     if ((count += 1) > 20) return false; // complexity limit
2974     for (uint i = 1; i < n->req(); i++) {
2975         Node* m = n->in(i);
2976         if (m == NULL || m == n || m->is_top()) continue;
2977         uint first_i = n->find_edge(m);
2978         if (i != first_i) continue; // process duplicate edge just once
2979         if (!detect_init_independence(m, st_is_pinned, count)) {
2980             return false;
2981         }
2982     }

2984     return true;
2985 }

2987 // Here are all the checks a Store must pass before it can be moved into
2988 // an initialization. Returns zero if a check fails.
2989 // On success, returns the (constant) offset to which the store applies,
2990 // within the initialized memory.
2991 intptr_t InitializeNode::can_capture_store(StoreNode* st, PhaseTransform* phase)
2992     const int FAIL = 0;

```

```

2993     if (st->req() != MemNode::ValueIn + 1)
2994         return FAIL; // an inscrutable StoreNode (card mark?)
2995     Node* ctl = st->in(MemNode::Control);
2996     if (!ctl != NULL && ctl->is_Proj() && ctl->in(0) == this)
2997         return FAIL; // must be unconditional after the initializatio
2998     Node* mem = st->in(MemNode::Memory);
2999     if (!mem->is_Proj() && mem->in(0) == this)
3000         return FAIL; // must not be preceded by other stores
3001     Node* adr = st->in(MemNode::Address);
3002     intptr_t offset;
3003     AllocateNode* alloc = AllocateNode::Ideal_allocation(adr, phase, offset);
3004     if (alloc == NULL)
3005         return FAIL; // inscrutable address
3006     if (alloc != allocation())
3007         return FAIL; // wrong allocation! (store needs to float up)
3008     Node* val = st->in(MemNode::ValueIn);
3009     int complexity_count = 0;
3010     if (!detect_init_independence(val, true, complexity_count))
3011         return FAIL; // stored value must be 'simple enough'

3013     return offset; // success
3014 }

3016 // Find the captured store in(i) which corresponds to the range
3017 // [start..start+size) in the initialized object.
3018 // If there is one, return its index i. If there isn't, return the
3019 // negative of the index where it should be inserted.
3020 // Return 0 if the queried range overlaps an initialization boundary
3021 // or if dead code is encountered.
3022 // If size_in_bytes is zero, do not bother with overlap checks.
3023 int InitializeNode::captured_store_insertion_point(intptr_t start,
3024     int size_in_bytes,
3025     PhaseTransform* phase) {
3026     const int FAIL = 0, MAX_STORE = BytesPerLong;

3028     if (is_complete())
3029         return FAIL; // arraycopy got here first; punt

3031     assert(allocation() != NULL, "must be present");

3033     // no negatives, no header fields:
3034     if (start < (intptr_t) allocation()->minimum_header_size()) return FAIL;

3036     // after a certain size, we bail out on tracking all the stores:
3037     intptr_t ti_limit = (TrackedInitializationLimit * HeapWordSize);
3038     if (start >= ti_limit) return FAIL;

3040     for (uint i = InitializeNode::RawStores, limit = req(); ; ) {
3041         if (i >= limit) return -(int)i; // not found; here is where to put it

3043         Node* st = in(i);
3044         intptr_t st_off = get_store_offset(st, phase);
3045         if (st_off < 0) {
3046             if (st != zero_memory()) {
3047                 return FAIL; // bail out if there is dead garbage
3048             }
3049         } else if (st_off > start) {
3050             // ...we are done, since stores are ordered
3051             if (st_off < start + size_in_bytes) {
3052                 return FAIL; // the next store overlaps
3053             }
3054             return -(int)i; // not found; here is where to put it
3055         } else if (st_off < start) {
3056             if (size_in_bytes != 0 &&
3057                 start < st_off + MAX_STORE &&
3058                 start < st_off + st->as_Store()->memory_size()) {

```

```

3059     return FAIL;          // the previous store overlaps
3060   }
3061   } else {
3062     if (size_in_bytes != 0 &&
3063         st->as_Store()->memory_size() != size_in_bytes) {
3064       return FAIL;      // mismatched store size
3065     }
3066     return i;
3067   }
3069   ++i;
3070 }
3071 }

3073 // Look for a captured store which initializes at the offset 'start'
3074 // with the given size. If there is no such store, and no other
3075 // initialization interferes, then return zero_memory (the memory
3076 // projection of the AllocateNode).
3077 Node* InitializeNode::find_captured_store(intptr_t start, int size_in_bytes,
3078                                         PhaseTransform* phase) {
3079   assert(stores_are_sane(phase), "");
3080   int i = captured_store_insertion_point(start, size_in_bytes, phase);
3081   if (i == 0) {
3082     return NULL;          // something is dead
3083   } else if (i < 0) {
3084     return zero_memory(); // just primordial zero bits here
3085   } else {
3086     Node* st = in(i);     // here is the store at this position
3087     assert(get_store_offset(st->as_Store(), phase) == start, "sanity");
3088     return st;
3089   }
3090 }

3092 // Create, as a raw pointer, an address within my new object at 'offset'.
3093 Node* InitializeNode::make_raw_address(intptr_t offset,
3094                                       PhaseTransform* phase) {
3095   Node* addr = in(RawAddress);
3096   if (offset != 0) {
3097     Compile* C = phase->C;
3098     addr = phase->transform( new (C, 4) AddPNode(C->top(), addr,
3099                                               phase->MakeConX(offset)) );
3100   }
3101   return addr;
3102 }

3104 // Clone the given store, converting it into a raw store
3105 // initializing a field or element of my new object.
3106 // Caller is responsible for retiring the original store,
3107 // with subsume_node or the like.
3108 //
3109 // From the example above InitializeNode::InitializeNode,
3110 // here are the old stores to be captured:
3111 //   store1 = (StoreC init.Control init.Memory (+ oop 12) 1)
3112 //   store2 = (StoreC init.Control store1      (+ oop 14) 2)
3113 //
3114 // Here is the changed code; note the extra edges on init:
3115 //   alloc = (Allocate ...)
3116 //   rawoop = alloc.RawAddress
3117 //   rawstore1 = (StoreC alloc.Control alloc.Memory (+ rawoop 12) 1)
3118 //   rawstore2 = (StoreC alloc.Control alloc.Memory (+ rawoop 14) 2)
3119 //   init = (Initialize alloc.Control alloc.Memory rawoop
3120 //          rawstore1 rawstore2)
3121 //
3122 Node* InitializeNode::capture_store(StoreNode* st, intptr_t start,
3123                                    PhaseTransform* phase) {
3124   assert(stores_are_sane(phase), "");

```

```

3126   if (start < 0) return NULL;
3127   assert(can_capture_store(st, phase) == start, "sanity");

3129   Compile* C = phase->C;
3130   int size_in_bytes = st->memory_size();
3131   int i = captured_store_insertion_point(start, size_in_bytes, phase);
3132   if (i == 0) return NULL; // bail out
3133   Node* prev_mem = NULL;   // raw memory for the captured store
3134   if (i > 0) {
3135     prev_mem = in(i);      // there is a pre-existing store under this one
3136     set_req(i, C->top());  // temporarily disconnect it
3137     // See StoreNode::Ideal 'st->outcnt() == 1' for the reason to disconnect.
3138   } else {
3139     i = -i;                // no pre-existing store
3140     prev_mem = zero_memory(); // a slice of the newly allocated object
3141     if (i > InitializeNode::RawStores && in(i-1) == prev_mem)
3142       set_req(--i, C->top()); // reuse this edge; it has been folded away
3143     else
3144       ins_req(i, C->top()); // build a new edge
3145   }
3146   Node* new_st = st->clone();
3147   new_st->set_req(MemNode::Control, in(Control));
3148   new_st->set_req(MemNode::Memory, prev_mem);
3149   new_st->set_req(MemNode::Address, make_raw_address(start, phase));
3150   new_st = phase->transform(new_st);

3152   // At this point, new_st might have swallowed a pre-existing store
3153   // at the same offset, or perhaps new_st might have disappeared,
3154   // if it redundantly stored the same value (or zero to fresh memory).

3156   // In any case, wire it in:
3157   set_req(i, new_st);

3159   // The caller may now kill the old guy.
3160   DEBUG_ONLY(Node* check_st = find_captured_store(start, size_in_bytes, phase));
3161   assert(check_st == new_st || check_st == NULL, "must be findable");
3162   assert(!is_complete(), "");
3163   return new_st;
3164 }

3166 static bool store_constant(jlong* tiles, int num_tiles,
3167                           intptr_t st_off, int st_size,
3168                           jlong con) {
3169   if ((st_off & (st_size-1)) != 0)
3170     return false; // strange store offset (assume size==2**N)
3171   address addr = (address)tiles + st_off;
3172   assert(st_off >= 0 && addr+st_size <= (address)&tiles[num_tiles], "oob");
3173   switch (st_size) {
3174     case sizeof(jbyte): *(jbyte*) addr = (jbyte) con; break;
3175     case sizeof(jchar): *(jchar*) addr = (jchar) con; break;
3176     case sizeof(jint): *(jint*) addr = (jint) con; break;
3177     case sizeof(jlong): *(jlong*) addr = (jlong) con; break;
3178     default: return false; // strange store size (detect size!=2**N here)
3179   }
3180   return true; // return success to caller
3181 }

3183 // Coalesce subword constants into int constants and possibly
3184 // into long constants. The goal, if the CPU permits,
3185 // is to initialize the object with a small number of 64-bit tiles.
3186 // Also, convert floating-point constants to bit patterns.
3187 // Non-constants are not relevant to this pass.
3188 //
3189 // In terms of the running example on InitializeNode::InitializeNode
3190 // and InitializeNode::capture_store, here is the transformation

```

```

3191 // of rawstore1 and rawstore2 into rawstore12:
3192 // alloc = (Allocate ...)
3193 // rawoop = alloc.RawAddress
3194 // tile12 = 0x00010002
3195 // rawstore12 = (StoreI alloc.Control alloc.Memory (+ rawoop 12) tile12)
3196 // init = (Initialize alloc.Control alloc.Memory rawoop rawstore12)
3197 //
3198 void
3199 InitializeNode::coalesce_subword_stores(intptr_t header_size,
3200                                         Node* size_in_bytes,
3201                                         PhaseGVN* phase) {
3202   Compile* C = phase->C;

3204   assert(stores_are_sane(phase), "");
3205   // Note: After this pass, they are not completely sane,
3206   // since there may be some overlaps.

3208   int old_subword = 0, old_long = 0, new_int = 0, new_long = 0;

3210   intptr_t ti_limit = (TrackedInitializationLimit * HeapWordSize);
3211   intptr_t size_limit = phase->find_intptr_t_con(size_in_bytes, ti_limit);
3212   size_limit = MIN2(size_limit, ti_limit);
3213   size_limit = align_size_up(size_limit, BytesPerLong);
3214   int num_tiles = size_limit / BytesPerLong;

3216   // allocate space for the tile map:
3217   const int small_len = DEBUG_ONLY(true ? 3 :) 30; // keep stack frames small
3218   jlong tiles_buf[small_len];
3219   Node* nodes_buf[small_len];
3220   jlong inits_buf[small_len];
3221   jlong* tiles = ((num_tiles <= small_len) ? &tiles_buf[0]
3222                  : NEW_RESOURCE_ARRAY(jlong, num_tiles));
3223   Node** nodes = ((num_tiles <= small_len) ? &nodes_buf[0]
3224                  : NEW_RESOURCE_ARRAY(Node*, num_tiles));
3225   jlong* inits = ((num_tiles <= small_len) ? &inits_buf[0]
3226                  : NEW_RESOURCE_ARRAY(jlong, num_tiles));
3227   // tiles: exact bitwise model of all primitive constants
3228   // nodes: last constant-storing node subsumed into the tiles model
3229   // inits: which bytes (in each tile) are touched by any initializations

3231   //// Pass A: Fill in the tile model with any relevant stores.

3233   Copy::zero_to_bytes(tiles, sizeof(tiles[0]) * num_tiles);
3234   Copy::zero_to_bytes(nodes, sizeof(nodes[0]) * num_tiles);
3235   Copy::zero_to_bytes(inits, sizeof(inits[0]) * num_tiles);
3236   Node* zmem = zero_memory(); // initially zero memory state
3237   for (uint i = InitializeNode::RawStores, limit = req(); i < limit; i++) {
3238     Node* st = in(i);
3239     intptr_t st_off = get_store_offset(st, phase);

3241     // Figure out the store's offset and constant value:
3242     if (st_off < header_size) continue; //skip (ignore header)
3243     if (st->in(MemNode::Memory) != zmem) continue; //skip (odd store chain)
3244     int st_size = st->as_Store()->memory_size();
3245     if (st_off + st_size > size_limit) break;

3247     // Record which bytes are touched, whether by constant or not.
3248     if (!store_constant(inits, num_tiles, st_off, st_size, (jlong) -1))
3249         continue; // skip (strange store size)

3251     const Type* val = phase->type(st->in(MemNode::ValueIn));
3252     if (!val->singleton()) continue; //skip (non-con store)
3253     BasicType type = val->basic_type();

3255     jlong con = 0;
3256     switch (type) {

```

```

3257     case T_INT: con = val->is_int()->get_con(); break;
3258     case T_LONG: con = val->is_long()->get_con(); break;
3259     case T_FLOAT: con = jint_cast(val->getf()); break;
3260     case T_DOUBLE: con = jlong_cast(val->getd()); break;
3261     default: continue; //skip (odd store type)
3262   }

3264   if (type == T_LONG && Matcher::isSimpleConstant64(con) &&
3265       st->Opcode() == Op_StoreL) {
3266     continue; // This StoreL is already optimal.
3267   }

3269   // Store down the constant.
3270   store_constant(tiles, num_tiles, st_off, st_size, con);

3272   intptr_t j = st_off >> LogBytesPerLong;

3274   if (type == T_INT && st_size == BytesPerInt
3275       && (st_off & BytesPerInt) == BytesPerInt) {
3276     jlong lcon = tiles[j];
3277     if (!Matcher::isSimpleConstant64(lcon) &&
3278         st->Opcode() == Op_StoreI) {
3279       // This StoreI is already optimal by itself.
3280       jint* intcon = (jint*) &tiles[j];
3281       intcon[1] = 0; // undo the store_constant()

3283       // If the previous store is also optimal by itself, back up and
3284       // undo the action of the previous loop iteration... if we can.
3285       // But if we can't, just let the previous half take care of itself.
3286       st = nodes[j];
3287       st_off -= BytesPerInt;
3288       con = intcon[0];
3289       if (con != 0 && st != NULL && st->Opcode() == Op_StoreI) {
3290         assert(st_off >= header_size, "still ignoring header");
3291         assert(get_store_offset(st, phase) == st_off, "must be");
3292         assert(in(i-1) == zmem, "must be");
3293         DEBUG_ONLY(const Type* tcon = phase->type(st->in(MemNode::ValueIn)));
3294         assert(con == tcon->is_int()->get_con(), "must be");
3295         // Undo the effects of the previous loop trip, which swallowed st:
3296         intcon[0] = 0; // undo store_constant()
3297         set_req(i-1, st); // undo set_req(i, zmem)
3298         nodes[j] = NULL; // undo nodes[j] = st
3299         --old_subword; // undo ++old_subword
3300       }
3301       continue; // This StoreI is already optimal.
3302     }
3303   }

3305   // This store is not needed.
3306   set_req(i, zmem);
3307   nodes[j] = st; // record for the moment
3308   if (st_size < BytesPerLong) // something has changed
3309     ++old_subword; // includes int/float, but who's counting...
3310   else ++old_long;
3311 }

3313 if ((old_subword + old_long) == 0)
3314   return; // nothing more to do

3316   //// Pass B: Convert any non-zero tiles into optimal constant stores.
3317   // Be sure to insert them before overlapping non-constant stores.
3318   // (E.g., byte[] x = { 1,2,y,4 } => x[int 0] = 0x01020004, x[2]=y.)
3319   for (int j = 0; j < num_tiles; j++) {
3320     jlong con = tiles[j];
3321     jlong init = inits[j];
3322     if (con == 0) continue;

```

```

3323 jint con0, con1;           // split the constant, address-wise
3324 jint init0, init1;       // split the init map, address-wise
3325 { union { jlong con; jint intcon[2]; } u;
3326   u.con = con;
3327   con0 = u.intcon[0];
3328   con1 = u.intcon[1];
3329   u.con = init;
3330   init0 = u.intcon[0];
3331   init1 = u.intcon[1];
3332 }

3334 Node* old = nodes[j];
3335 assert(old != NULL, "need the prior store");
3336 intptr_t offset = (j * BytesPerLong);

3338 bool split = !Matcher::isSimpleConstant64(con);

3340 if (offset < header_size) {
3341   assert(offset + BytesPerInt >= header_size, "second int counts");
3342   assert(*(jint*)&tiles[j] == 0, "junk in header");
3343   split = true;           // only the second word counts
3344   // Example: int a[] = { 42 ... }
3345 } else if (con0 == 0 && init0 == -1) {
3346   split = true;         // first word is covered by full inits
3347   // Example: int a[] = { ... foo(), 42 ... }
3348 } else if (con1 == 0 && init1 == -1) {
3349   split = true;         // second word is covered by full inits
3350   // Example: int a[] = { ... 42, foo() ... }
3351 }

3353 // Here's a case where init0 is neither 0 nor -1:
3354 // byte a[] = { ... 0,0,foo(),0, 0,0,0,42 ... }
3355 // Assuming big-endian memory, init0, init1 are 0x0000FF00, 0x000000FF.
3356 // In this case the tile is not split; it is (jlong)42.
3357 // The big tile is stored down, and then the foo(),0 value is inserted.
3358 // (If there were foo(),foo() instead of foo(),0, init0 would be -1.)

3360 Node* ctl = old->in(MemNode::Control);
3361 Node* adr = make_raw_address(offset, phase);
3362 const TypePtr* atp = TypeRawPtr::BOTTOM;

3364 // One or two coalesced stores to plop down.
3365 Node* st[2];
3366 intptr_t off[2];
3367 int nst = 0;
3368 if (!split) {
3369   ++new_long;
3370   off[nst] = offset;
3371   st[nst++] = StoreNode::make(*phase, ctl, zmem, adr, atp,
3372                               phase->longcon(con), T_LONG);
3373 } else {
3374   // Omit either if it is a zero.
3375   if (con0 != 0) {
3376     ++new_int;
3377     off[nst] = offset;
3378     st[nst++] = StoreNode::make(*phase, ctl, zmem, adr, atp,
3379                               phase->intcon(con0), T_INT);
3380   }
3381   if (con1 != 0) {
3382     ++new_int;
3383     offset += BytesPerInt;
3384     adr = make_raw_address(offset, phase);
3385     off[nst] = offset;
3386     st[nst++] = StoreNode::make(*phase, ctl, zmem, adr, atp,
3387                               phase->intcon(con1), T_INT);
3388   }

```

```

3389   }

3391   // Insert second store first, then the first before the second.
3392   // Insert each one just before any overlapping non-constant stores.
3393   while (nst > 0) {
3394     Node* st1 = st[--nst];
3395     C->copy_node_notes_to(st1, old);
3396     st1 = phase->transform(st1);
3397     offset = off[nst];
3398     assert(offset >= header_size, "do not smash header");
3399     int ins_idx = captured_store_insertion_point(offset, /*size:*/0, phase);
3400     guarantee(ins_idx != 0, "must re-insert constant store");
3401     if (ins_idx < 0) ins_idx = -ins_idx; // never overlap
3402     if (ins_idx > InitializeNode::RawStores && in(ins_idx-1) == zmem)
3403       set_req(--ins_idx, st1);
3404     else
3405       ins_req(ins_idx, st1);
3406   }
3407 }

3409 if (PrintCompilation && WizardMode)
3410   tty->print_cr("Changed %d/%d subword/long constants into %d/%d int/long",
3411               old_subword, old_long, new_int, new_long);
3412 if (C->log() != NULL)
3413   C->log()->elem("comment that='%d/%d subword/long to %d/%d int/long'",
3414               old_subword, old_long, new_int, new_long);

3416 // Clean up any remaining occurrences of zmem:
3417 remove_extra zeroes();
3418 }

3420 // Explore forward from in(start) to find the first fully initialized
3421 // word, and return its offset. Skip groups of subword stores which
3422 // together initialize full words. If in(start) is itself part of a
3423 // fully initialized word, return the offset of in(start). If there
3424 // are no following full-word stores, or if something is fishy, return
3425 // a negative value.
3426 intptr_t InitializeNode::find_next_fullword_store(uint start, PhaseGVN* phase) {
3427   int int_map = 0;
3428   intptr_t int_map_off = 0;
3429   const int FULL_MAP = right_n_bits(BytesPerInt); // the int_map we hope for

3431   for (uint i = start, limit = req(); i < limit; i++) {
3432     Node* st = in(i);

3434     intptr_t st_off = get_store_offset(st, phase);
3435     if (st_off < 0) break; // return conservative answer

3437     int st_size = st->as_Store()->memory_size();
3438     if (st_size >= BytesPerInt && (st_off % BytesPerInt) == 0) {
3439       return st_off; // we found a complete word init
3440     }

3442     // update the map:

3444     intptr_t this_int_off = align_size_down(st_off, BytesPerInt);
3445     if (this_int_off != int_map_off) {
3446       // reset the map:
3447       int_map = 0;
3448       int_map_off = this_int_off;
3449     }

3451     int subword_off = st_off - this_int_off;
3452     int_map |= right_n_bits(st_size) << subword_off;
3453     if ((int_map & FULL_MAP) == FULL_MAP) {
3454       return this_int_off; // we found a complete word init

```

```

3455     }
3457     // Did this store hit or cross the word boundary?
3458     intptr_t next_int_off = align_size_down(st_off + st_size, BytesPerInt);
3459     if (next_int_off == this_int_off + BytesPerInt) {
3460         // We passed the current int, without fully initializing it.
3461         int_map_off = next_int_off;
3462         int_map >>= BytesPerInt;
3463     } else if (next_int_off > this_int_off + BytesPerInt) {
3464         // We passed the current and next int.
3465         return this_int_off + BytesPerInt;
3466     }
3467 }
3469 return -1;
3470 }

3473 // Called when the associated AllocateNode is expanded into CFG.
3474 // At this point, we may perform additional optimizations.
3475 // Linearize the stores by ascending offset, to make memory
3476 // activity as coherent as possible.
3477 Node* InitializeNode::complete_stores(Node* rawctl, Node* rawmem, Node* rawptr,
3478                                     intptr_t header_size,
3479                                     Node* size_in_bytes,
3480                                     PhaseGVN* phase) {
3481     assert(!is_complete(), "not already complete");
3482     assert(stores_are_sane(phase), "");
3483     assert(allocation() != NULL, "must be present");

3485     remove_extra_zeroes();

3487     if (ReduceFieldZeroing || ReduceBulkZeroing)
3488         // reduce instruction count for common initialization patterns
3489         coalesce_subword_stores(header_size, size_in_bytes, phase);

3491     Node* zmem = zero_memory(); // initially zero memory state
3492     Node* inits = zmem;        // accumulating a linearized chain of inits
3493     #ifdef ASSERT
3494     intptr_t first_offset = allocation()->minimum_header_size();
3495     intptr_t last_init_off = first_offset; // previous init offset
3496     intptr_t last_init_end = first_offset; // previous init offset+size
3497     intptr_t last_tile_end = first_offset; // previous tile offset+size
3498     #endif
3499     intptr_t zeroes_done = header_size;

3501     bool do_zeroing = true; // we might give up if inits are very sparse
3502     int big_init_gaps = 0; // how many large gaps have we seen?

3504     if (ZeroTLAB) do_zeroing = false;
3505     if (!ReduceFieldZeroing && !ReduceBulkZeroing) do_zeroing = false;

3507     for (uint i = InitializeNode::RawStores, limit = req(); i < limit; i++) {
3508         Node* st = in(i);
3509         intptr_t st_off = get_store_offset(st, phase);
3510         if (st_off < 0)
3511             break; // unknown junk in the inits
3512         if (st->in(MemNode::Memory) != zmem)
3513             break; // complicated store chains somehow in list

3515         int st_size = st->as_Store()->memory_size();
3516         intptr_t next_init_off = st_off + st_size;

3518         if (do_zeroing && zeroes_done < next_init_off) {
3519             // See if this store needs a zero before it or under it.
3520             intptr_t zeroes_needed = st_off;

```

```

3522         if (st_size < BytesPerInt) {
3523             // Look for subword stores which only partially initialize words.
3524             // If we find some, we must lay down some word-level zeroes first,
3525             // underneath the subword stores.
3526             //
3527             // Examples:
3528             //   byte[] a = { p,q,r,s } => a[0]=p,a[1]=q,a[2]=r,a[3]=s
3529             //   byte[] a = { x,y,0,0 } => a[0..3] = 0, a[0]=x,a[1]=y
3530             //   byte[] a = { 0,0,z,0 } => a[0..3] = 0, a[2]=z
3531             //
3532             // Note: coalesce_subword_stores may have already done this,
3533             // if it was prompted by constant non-zero subword initializers.
3534             // But this case can still arise with non-constant stores.

3536             intptr_t next_full_store = find_next_fullword_store(i, phase);

3538             // In the examples above:
3539             //   in(i)           p  q  r  s      x  y  z
3540             //   st_off         12 13 14 15     12 13 14
3541             //   st_size        1  1  1  1      1  1  1
3542             //   next_full_s.   12 16 16 16     16 16 16
3543             //   z's_done       12 16 16 16     12 16 12
3544             //   z's_needed     12 16 16 16     16 16 16
3545             //   zsize          0  0  0  0      4  0  4
3546             if (next_full_store < 0) {
3547                 // Conservative tack: Zero to end of current word.
3548                 zeroes_needed = align_size_up(zeroes_needed, BytesPerInt);
3549             } else {
3550                 // Zero to beginning of next fully initialized word.
3551                 // Or, don't zero at all, if we are already in that word.
3552                 assert(next_full_store >= zeroes_needed, "must go forward");
3553                 assert((next_full_store & (BytesPerInt-1)) == 0, "even boundary");
3554                 zeroes_needed = next_full_store;
3555             }
3556         }

3558         if (zeroes_needed > zeroes_done) {
3559             intptr_t zsize = zeroes_needed - zeroes_done;
3560             // Do some incremental zeroing on rawmem, in parallel with inits.
3561             zeroes_done = align_size_down(zeroes_done, BytesPerInt);
3562             rawmem = ClearArrayNode::clear_memory(rawctl, rawmem, rawptr,
3563                                                  zeroes_done, zeroes_needed,
3564                                                  phase);
3565             zeroes_done = zeroes_needed;
3566             if (zsize > Matcher::init_array_short_size && ++big_init_gaps > 2)
3567                 do_zeroing = false; // leave the hole, next time
3568         }
3569     }

3571     // Collect the store and move on:
3572     st->set_req(MemNode::Memory, inits);
3573     inits = st; // put it on the linearized chain
3574     set_req(i, zmem); // unhook from previous position

3576     if (zeroes_done == st_off)
3577         zeroes_done = next_init_off;

3579     assert(!do_zeroing || zeroes_done >= next_init_off, "don't miss any");

3581     #ifdef ASSERT
3582     // Various order invariants. Weaker than stores_are_sane because
3583     // a large constant tile can be filled in by smaller non-constant stores.
3584     assert(st_off >= last_init_off, "inits do not reverse");
3585     last_init_off = st_off;
3586     const Type* val = NULL;

```



```

3587     if (st_size >= BytesPerInt &&
3588         (val = phase->type(st->in(MemNode::ValueIn))>singleton() &&
3589          (int)val->basic_type() < (int)T_OBJECT) {
3590         assert(st_off >= last_tile_end, "tiles do not overlap");
3591         assert(st_off >= last_init_end, "tiles do not overwrite inits");
3592         last_tile_end = MAX2(last_tile_end, next_init_off);
3593     } else {
3594         intptr_t st_tile_end = align_size_up(next_init_off, BytesPerLong);
3595         assert(st_tile_end >= last_tile_end, "inits stay with tiles");
3596         assert(st_off >= last_init_end, "inits do not overlap");
3597         last_init_end = next_init_off; // it's a non-tile
3598     }
3599 }
3600 #endif //ASSERT

3602 remove_extra_zeroes(); // clear out all the zmems left over
3603 add_req(inits);

3605 if (!ZeroTLAB) {
3606     // If anything remains to be zeroed, zero it all now.
3607     zeroes_done = align_size_down(zeroes_done, BytesPerInt);
3608     // if it is the last unused 4 bytes of an instance, forget about it
3609     intptr_t size_limit = phase->find_intptr_t_con(size_in_bytes, max_jint);
3610     if (zeroes_done + BytesPerLong >= size_limit) {
3611         assert(allocation() != NULL, "");
3612         if (allocation()->Opcode() == Op_Allocate) {
3613             Node* klass_node = allocation()->in(AllocateNode::KlassNode);
3614             cKlass* k = phase->type(klass_node)->is_klassptr()->klass();
3615             if (zeroes_done == k->layout_helper())
3616                 zeroes_done = size_limit;
3617         }
3618     }
3619     if (zeroes_done < size_limit) {
3620         rawmem = ClearArrayNode::clear_memory(rawctl, rawmem, rawptr,
3621         zeroes_done, size_in_bytes, phase);
3622     }
3623 }

3625 set_complete(phase);
3626 return rawmem;
3627 }

3630 #ifndef ASSERT
3631 bool InitializeNode::stores_are_sane(PhaseTransform* phase) {
3632     if (is_complete())
3633         return true; // stores could be anything at this point
3634     assert(allocation() != NULL, "must be present");
3635     intptr_t last_off = allocation()->minimum_header_size();
3636     for (uint i = InitializeNode::RawStores; i < req(); i++) {
3637         Node* st = in(i);
3638         intptr_t st_off = get_store_offset(st, phase);
3639         if (st_off < 0) continue; // ignore dead garbage
3640         if (last_off > st_off) {
3641             tty->print_cr("*** bad store offset at %d: %d > %d", i, last_off, st_off);
3642             this->dump(2);
3643             assert(false, "ascending store offsets");
3644             return false;
3645         }
3646         last_off = st_off + st->as_Store()->memory_size();
3647     }
3648     return true;
3649 }
3650 #endif //ASSERT

```

```

3655 //=====MergeMemNode=====
3656 //
3657 // SEMANTICS OF MEMORY MERGES: A MergeMem is a memory state assembled from several
3658 // contributing store or call operations. Each contributor provides the memory
3659 // state for a particular "alias type" (see Compile::alias_type). For example,
3660 // if a MergeMem has an input X for alias category #6, then any memory reference
3661 // to alias category #6 may use X as its memory state input, as an exact equivalent
3662 // to using the MergeMem as a whole.
3663 // Load<6>( MergeMem(<6>: X, ...), p ) <==> Load<6>(X,p)
3664 //
3665 // (Here, the <N> notation gives the index of the relevant adr_type.)
3666 //
3667 // In one special case (and more cases in the future), alias categories overlap.
3668 // The special alias category "Bot" (Compile::AliasIdxBot) includes all memory
3669 // states. Therefore, if a MergeMem has only one contributing input W for Bot,
3670 // it is exactly equivalent to that state W:
3671 // MergeMem(<Bot>: W) <==> W
3672 //
3673 // Usually, the merge has more than one input. In that case, where inputs
3674 // overlap (i.e., one is Bot), the narrower alias type determines the memory
3675 // state for that type, and the wider alias type (Bot) fills in everywhere else:
3676 // Load<5>( MergeMem(<Bot>: W, <6>: X), p ) <==> Load<5>(W,p)
3677 // Load<6>( MergeMem(<Bot>: W, <6>: X), p ) <==> Load<6>(X,p)
3678 //
3679 // A merge can take a "wide" memory state as one of its narrow inputs.
3680 // This simply means that the merge observes only the relevant parts of
3681 // the wide input. That is, wide memory states arriving at narrow merge inputs
3682 // are implicitly "filtered" or "sliced" as necessary. (This is rare.)
3683 //
3684 // These rules imply that MergeMem nodes may cascade (via their <Bot> links),
3685 // and that memory slices "leak through":
3686 // MergeMem(<Bot>: MergeMem(<Bot>: W, <7>: Y)) <==> MergeMem(<Bot>: W, <7>: Y)
3687 //
3688 // But, in such a cascade, repeated memory slices can "block the leak":
3689 // MergeMem(<Bot>: MergeMem(<Bot>: W, <7>: Y), <7>: Y') <==> MergeMem(<Bot>: W
3690 //
3691 // In the last example, Y is not part of the combined memory state of the
3692 // outermost MergeMem. The system must, of course, prevent unschedulable
3693 // memory states from arising, so you can be sure that the state Y is somehow
3694 // a precursor to state Y'.
3695 //
3696 //
3697 // REPRESENTATION OF MEMORY MERGES: The indexes used to address the Node::in arr
3698 // of each MergeMemNode array are exactly the numerical alias indexes, including
3699 // but not limited to AliasIdxTop, AliasIdxBot, and AliasIdxRaw. The functions
3700 // Compile::alias_type (and kin) produce and manage these indexes.
3701 //
3702 // By convention, the value of in(AliasIdxTop) (i.e., in(1)) is always the top n
3703 // (Note that this provides quick access to the top node inside MergeMem methods
3704 // without the need to reach out via TLS to Compile::current.)
3705 //
3706 // As a consequence of what was just described, a MergeMem that represents a full
3707 // memory state has an edge in(AliasIdxBot) which is a "wide" memory state,
3708 // containing all alias categories.
3709 //
3710 // MergeMem nodes never (?) have control inputs, so in(0) is NULL.
3711 //
3712 // All other edges in(N) (including in(AliasIdxRaw), which is in(3)) are either
3713 // a memory state for the alias type <N>, or else the top node, meaning that
3714 // there is no particular input for that alias type. Note that the length of
3715 // a MergeMem is variable, and may be extended at any time to accommodate new
3716 // memory states at larger alias indexes. When merges grow, they are of course
3717 // filled with "top" in the unused in() positions.
3718 //

```

```

3719 // This use of top is named "empty_memory()", or "empty_mem" (no-memory) as a va
3720 // (Top was chosen because it works smoothly with passes like GCM.)
3721 //
3722 // For convenience, we hardwire the alias index for TypeRawPtr::BOTTOM. (It is
3723 // the type of random VM bits like TLS references.) Since it is always the
3724 // first non-Bot memory slice, some low-level loops use it to initialize an
3725 // index variable: for (i = AliasIdxRaw; i < req(); i++).
3726 //
3727 //
3728 // ACCESSORS: There is a special accessor MergeMemNode::base_memory which retur
3729 // the distinguished "wide" state. The accessor MergeMemNode::memory_at(N) retu
3730 // the memory state for alias type <N>, or (if there is no particular slice at <
3731 // it returns the base memory. To prevent bugs, memory_at does not accept <Top>
3732 // or <Bot> indexes. The iterator MergeMemStream provides robust iteration over
3733 // MergeMem nodes or pairs of such nodes, ensuring that the non-top edges are vi
3734 //
3735 // %%% We may get rid of base_memory as a separate accessor at some point; it i
3736 // really that different from the other memory inputs. An abbreviation called
3737 // "bot_memory()" for "memory_at(AliasIdxBot)" would keep code tidy.
3738 //
3739 //
3740 // PARTIAL MEMORY STATES: During optimization, MergeMem nodes may arise that re
3741 // partial memory states. When a Phi splits through a MergeMem, the copy of the
3742 // that "emerges though" the base memory will be marked as excluding the alias t
3743 // of the other (narrow-memory) copies which "emerged through" the narrow edges:
3744 //
3745 // Phi<Bot>(U, MergeMem(<Bot>: W, <8>: Y))
3746 // ==Ideal=> MergeMem(<Bot>: Phi<Bot-8>(U, W), Phi<8>(U, Y))
3747 //
3748 // This strange "subtraction" effect is necessary to ensure IGVN convergence.
3749 // (It is currently unimplemented.) As you can see, the resulting merge is
3750 // actually a disjoint union of memory states, rather than an overlay.
3751 //
3752 //
3753 //-----MergeMemNode-----
3754 Node* MergeMemNode::make_empty_memory() {
3755     Node* empty_memory = (Node*) Compile::current()->top();
3756     assert(empty_memory->is_top(), "correct sentinel identity");
3757     return empty_memory;
3758 }
3759
3760 MergeMemNode::MergeMemNode(Node *new_base) : Node(1+Compile::AliasIdxRaw) {
3761     init_class_id(Class_MergeMem);
3762     // all inputs are nullified in Node::Node(int)
3763     // set_input(0, NULL); // no control input
3764
3765     // Initialize the edges uniformly to top, for starters.
3766     Node* empty_mem = make_empty_memory();
3767     for (uint i = Compile::AliasIdxTop; i < req(); i++) {
3768         init_req(i, empty_mem);
3769     }
3770     assert(empty_memory() == empty_mem, "");
3771
3772     if (new_base != NULL && new_base->is_MergeMem()) {
3773         MergeMemNode* mdef = new_base->as_MergeMem();
3774         assert(mdef->empty_memory() == empty_mem, "consistent sentinels");
3775         for (MergeMemStream mms(this, mdef); mms.next_non_empty2(); ) {
3776             mms.set_memory(mms.memory2());
3777         }
3778         assert(base_memory() == mdef->base_memory(), "");
3779     } else {
3780         set_base_memory(new_base);
3781     }
3782 }
3783
3784 // Make a new, untransformed MergeMem with the same base as 'mem'.

```

```

3785 // If mem is itself a MergeMem, populate the result with the same edges.
3786 MergeMemNode* MergeMemNode::make(Compile* C, Node* mem) {
3787     return new(C, 1+Compile::AliasIdxRaw) MergeMemNode(mem);
3788 }
3789
3790 //-----cmp-----
3791 uint MergeMemNode::hash() const { return NO_HASH; }
3792 uint MergeMemNode::cmp(const Node &n) const {
3793     return (&n == this); // Always fail except on self
3794 }
3795
3796 //-----Identity-----
3797 Node* MergeMemNode::Identity(PhaseTransform *phase) {
3798     // Identity if this merge point does not record any interesting memory
3799     // disambiguations.
3800     Node* base_mem = base_memory();
3801     Node* empty_mem = empty_memory();
3802     if (base_mem != empty_mem) { // Memory path is not dead?
3803         for (uint i = Compile::AliasIdxRaw; i < req(); i++) {
3804             Node* mem = in(i);
3805             if (mem != empty_mem && mem != base_mem) {
3806                 return this; // Many memory splits; no change
3807             }
3808         }
3809     }
3810     return base_mem; // No memory splits; ID on the one true input
3811 }
3812
3813 //-----Ideal-----
3814 // This method is invoked recursively on chains of MergeMem nodes
3815 Node *MergeMemNode::Ideal(PhaseGVN *phase, bool can_reshape) {
3816     // Remove chain'd MergeMems
3817     //
3818     // This is delicate, because the each "in(i)" (i >= Raw) is interpreted
3819     // relative to the "in(Bot)". Since we are patching both at the same time,
3820     // we have to be careful to read each "in(i)" relative to the old "in(Bot)",
3821     // but rewrite each "in(i)" relative to the new "in(Bot)".
3822     Node *progress = NULL;
3823
3824     Node* old_base = base_memory();
3825     Node* empty_mem = empty_memory();
3826     if (old_base == empty_mem)
3827         return NULL; // Dead memory path.
3828
3829     MergeMemNode* old_mbase;
3830     if (old_base != NULL && old_base->is_MergeMem())
3831         old_mbase = old_base->as_MergeMem();
3832     else
3833         old_mbase = NULL;
3834     Node* new_base = old_base;
3835
3836     // simplify stacked MergeMems in base memory
3837     if (old_mbase) new_base = old_mbase->base_memory();
3838
3839     // the base memory might contribute new slices beyond my req()
3840     if (old_mbase) grow_to_match(old_mbase);
3841
3842     // Look carefully at the base node if it is a phi.
3843     PhiNode* phi_base;
3844     if (new_base != NULL && new_base->is_Phi())
3845         phi_base = new_base->as_Phi();
3846     else
3847         phi_base = NULL;
3848
3849     Node* phi_reg = NULL;

```

```

3851 uint phi_len = (uint)-1;
3852 if (phi_base != NULL && !phi_base->is_copy()) {
3853     // do not examine phi if degraded to a copy
3854     phi_reg = phi_base->region();
3855     phi_len = phi_base->req();
3856     // see if the phi is unfinished
3857     for (uint i = 1; i < phi_len; i++) {
3858         if (phi_base->in(i) == NULL) {
3859             // incomplete phi; do not look at it yet!
3860             phi_reg = NULL;
3861             phi_len = (uint)-1;
3862             break;
3863         }
3864     }
3865 }

3867 // Note: We do not call verify_sparse on entry, because inputs
3868 // can normalize to the base memory via subsume_node or similar
3869 // mechanisms. This method repairs that damage.

3871 assert(!old_mbase || old_mbase->is_empty_memory(empty_mem), "consistent sentin

3873 // Look at each slice.
3874 for (uint i = Compile::AliasIdxRaw; i < req(); i++) {
3875     Node* old_in = in(i);
3876     // calculate the old memory value
3877     Node* old_mem = old_in;
3878     if (old_mem == empty_mem) old_mem = old_base;
3879     assert(old_mem == memory_at(i), "");

3881 // maybe update (reslice) the old memory value

3883 // simplify stacked MergeMems
3884 Node* new_mem = old_mem;
3885 MergeMemNode* old_mmem;
3886 if (old_mem != NULL && old_mem->is_MergeMem())
3887     old_mmem = old_mem->as_MergeMem();
3888 else
3889     old_mmem = NULL;
3890 if (old_mmem == this) {
3891     // This can happen if loops break up and safepoints disappear.
3892     // A merge of BotPtr (default) with a RawPtr memory derived from a
3893     // safepoint can be rewritten to a merge of the same BotPtr with
3894     // the BotPtr phi coming into the loop. If that phi disappears
3895     // also, we can end up with a self-loop of the mergemem.
3896     // In general, if loops degenerate and memory effects disappear,
3897     // a mergemem can be left looking at itself. This simply means
3898     // that the mergemem's default should be used, since there is
3899     // no longer any apparent effect on this slice.
3900     // Note: If a memory slice is a MergeMem cycle, it is unreachable
3901     // from start. Update the input to TOP.
3902     new_mem = (new_base == this || new_base == empty_mem)? empty_mem : new_base;
3903 }
3904 else if (old_mmem != NULL) {
3905     new_mem = old_mmem->memory_at(i);
3906 }
3907 // else preceding memory was not a MergeMem

3909 // replace equivalent this (unfortunately, they do not GVN together)
3910 if (new_mem != NULL && new_mem != new_base &&
3911     new_mem->req() == phi_len && new_mem->in(0) == phi_reg) {
3912     if (new_mem->is_Phi()) {
3913         PhiNode* phi_mem = new_mem->as_Phi();
3914         for (uint i = 1; i < phi_len; i++) {
3915             if (phi_base->in(i) != phi_mem->in(i)) {
3916                 phi_mem = NULL;

```

```

3917         break;
3918     }
3919     }
3920     if (phi_mem != NULL) {
3921         // equivalent phi nodes; revert to the def
3922         new_mem = new_base;
3923     }
3924 }
3925 }

3927 // maybe store down a new value
3928 Node* new_in = new_mem;
3929 if (new_in == new_base) new_in = empty_mem;

3931 if (new_in != old_in) {
3932     // Warning: Do not combine this "if" with the previous "if"
3933     // A memory slice might have be be rewritten even if it is semantically
3934     // unchanged, if the base_memory value has changed.
3935     set_req(i, new_in);
3936     progress = this; // Report progress
3937 }
3938 }

3940 if (new_base != old_base) {
3941     set_req(Compile::AliasIdxBot, new_base);
3942     // Don't use set_base_memory(new_base), because we need to update du.
3943     assert(base_memory() == new_base, "");
3944     progress = this;
3945 }

3947 if( base_memory() == this ) {
3948     // a self cycle indicates this memory path is dead
3949     set_req(Compile::AliasIdxBot, empty_mem);
3950 }

3952 // Resolve external cycles by calling Ideal on a MergeMem base_memory
3953 // Recursion must occur after the self cycle check above
3954 if( base_memory()->is_MergeMem() ) {
3955     MergeMemNode *new_mbase = base_memory()->as_MergeMem();
3956     Node *m = phase->transform(new_mbase); // Rollup any cycles
3957     if( m != NULL && (m->is_top() ||
3958         m->is_MergeMem() && m->as_MergeMem()->base_memory() == empty_mem) ) {
3959         // propagate rollup of dead cycle to self
3960         set_req(Compile::AliasIdxBot, empty_mem);
3961     }
3962 }

3964 if( base_memory() == empty_mem ) {
3965     progress = this;
3966     // Cut inputs during Parse phase only.
3967     // During Optimize phase a dead MergeMem node will be subsumed by Top.
3968     if( !can_reshape ) {
3969         for (uint i = Compile::AliasIdxRaw; i < req(); i++) {
3970             if (in(i) != empty_mem) { set_req(i, empty_mem); }
3971         }
3972     }
3973 }

3975 if( !progress && base_memory()->is_Phi() && can_reshape ) {
3976     // Check if PhiNode::Ideal's "Splitphis through memory merges"
3977     // transform should be attempted. Look for this->phi->this cycle.
3978     uint merge_width = req();
3979     if (merge_width > Compile::AliasIdxRaw) {
3980         PhiNode* phi = base_memory()->as_Phi();
3981         for( uint i = 1; i < phi->req(); ++i ) { // For all paths in
3982             if (phi->in(i) == this) {

```

```

3983     phase->is_IterGVN()->_worklist.push(phi);
3984     break;
3985   }
3986 }
3987 }
3988 }

3990 assert(progress || verify_sparse(), "please, no dups of base");
3991 return progress;
3992 }

3994 //-----set_base_memory-----
3995 void MergeMemNode::set_base_memory(Node *new_base) {
3996     Node* empty_mem = empty_memory();
3997     set_req(Compile::AliasIdxBot, new_base);
3998     assert(memory_at(req()) == new_base, "must set default memory");
3999     // Clear out other occurrences of new_base:
4000     if (new_base != empty_mem) {
4001         for (uint i = Compile::AliasIdxRaw; i < req(); i++) {
4002             if (in(i) == new_base) set_req(i, empty_mem);
4003         }
4004     }
4005 }

4007 //-----out_RegMask-----
4008 const RegMask &MergeMemNode::out_RegMask() const {
4009     return RegMask::Empty;
4010 }

4012 //-----dump_spec-----
4013 #ifndef PRODUCT
4014 void MergeMemNode::dump_spec(outputStream *st) const {
4015     st->print(" ");
4016     Node* base_mem = base_memory();
4017     for (uint i = Compile::AliasIdxRaw; i < req(); i++) {
4018         Node* mem = memory_at(i);
4019         if (mem == base_mem) { st->print(" -"); continue; }
4020         st->print(" N%d:", mem->idx );
4021         Compile::current()->get_adr_type(i)->dump_on(st);
4022     }
4023     st->print(" ");
4024 }
4025 #endif // !PRODUCT

4028 #ifdef ASSERT
4029 static bool might_be_same(Node* a, Node* b) {
4030     if (a == b) return true;
4031     if (!(a->is_Phi() || b->is_Phi())) return false;
4032     // this shift around during optimization
4033     return true; // pretty stupid...
4034 }

4036 // verify a narrow slice (either incoming or outgoing)
4037 static void verify_memory_slice(const MergeMemNode* m, int alias_idx, Node* n) {
4038     if (!VerifyAliases) return; // don't bother to verify unless requested
4039     if (is_error_reported()) return; // muzzle asserts when debugging an error
4040     if (Node::in_dump()) return; // muzzle asserts when printing
4041     assert(alias_idx >= Compile::AliasIdxRaw, "must not disturb base_memory or sen
4042     assert(n != NULL, "");
4043     // Elide intervening MergeMem's
4044     while (n->is_MergeMem()) {
4045         n = n->as_MergeMem()->memory_at(alias_idx);
4046     }
4047     Compile* C = Compile::current();
4048     const TypePtr* n_adr_type = n->adr_type();

```

```

4049     if (n == m->empty_memory()) {
4050         // Implicit copy of base_memory()
4051     } else if (n_adr_type != TypePtr::BOTTOM) {
4052         assert(n_adr_type != NULL, "new memory must have a well-defined adr_type");
4053         assert(C->must_alias(n_adr_type, alias_idx), "new memory must match selected
4054     } else {
4055         // A few places like make_runtime_call "know" that VM calls are narrow,
4056         // and can be used to update only the VM bits stored as TypeRawPtr::BOTTOM.
4057         bool expected_wide_mem = false;
4058         if (n == m->base_memory()) {
4059             expected_wide_mem = true;
4060         } else if (alias_idx == Compile::AliasIdxRaw ||
4061                 n == m->memory_at(Compile::AliasIdxRaw)) {
4062             expected_wide_mem = true;
4063         } else if (!C->alias_type(alias_idx)->is_rewritable()) {
4064             // memory can "leak through" calls on channels that
4065             // are write-once. Allow this also.
4066             expected_wide_mem = true;
4067         }
4068         assert(expected_wide_mem, "expected narrow slice replacement");
4069     }
4070 }
4071 #else // !ASSERT
4072 #define verify_memory_slice(m,i,n) (0) // PRODUCT version is no-op
4073 #endif

4076 //-----memory_at-----
4077 Node* MergeMemNode::memory_at(uint alias_idx) const {
4078     assert(alias_idx >= Compile::AliasIdxRaw ||
4079            alias_idx == Compile::AliasIdxBot && Compile::current()->AliasLevel() =
4080            "must avoid base_memory and AliasIdxTop");

4082     // Otherwise, it is a narrow slice.
4083     Node* n = alias_idx < req() ? in(alias_idx) : empty_memory();
4084     Compile *C = Compile::current();
4085     if (is_empty_memory(n)) {
4086         // the array is sparse; empty slots are the "top" node
4087         n = base_memory();
4088         assert(Node::in_dump()
4089                || n == NULL || n->bottom_type() == Type::TOP
4090                || n->adr_type() == NULL // address is TOP
4091                || n->adr_type() == TypePtr::BOTTOM
4092                || n->adr_type() == TypeRawPtr::BOTTOM
4093                || Compile::current()->AliasLevel() == 0,
4094                "must be a wide memory");
4095         // AliasLevel == 0 if we are organizing the memory states manually.
4096         // See verify_memory_slice for comments on TypeRawPtr::BOTTOM.
4097     } else {
4098         // make sure the stored slice is sane
4099         #ifdef ASSERT
4100             if (is_error_reported() || Node::in_dump()) {
4101                 } else if (might_be_same(n, base_memory())) {
4102                     // Give it a pass: It is a mostly harmless repetition of the base.
4103                     // This can arise normally from node subsumption during optimization.
4104                 } else {
4105                     verify_memory_slice(this, alias_idx, n);
4106                 }
4107             #endif
4108         }
4109         return n;
4110     }

4112 //-----set_memory_at-----
4113 void MergeMemNode::set_memory_at(uint alias_idx, Node *n) {
4114     verify_memory_slice(this, alias_idx, n);

```

```

4115 Node* empty_mem = empty_memory();
4116 if (n == base_memory()) n = empty_mem; // collapse default
4117 uint need_req = alias_idx+1;
4118 if (req() < need_req) {
4119     if (n == empty_mem) return; // already the default, so do not grow me
4120     // grow the sparse array
4121     do {
4122         add_req(empty_mem);
4123     } while (req() < need_req);
4124 }
4125 set_req( alias_idx, n );
4126 }

4130 //-----iteration_setup-----
4131 void MergeMemNode::iteration_setup(const MergeMemNode* other) {
4132     if (other != NULL) {
4133         grow_to_match(other);
4134         // invariant: the finite support of mm2 is within mm->req()
4135         #ifdef ASSERT
4136         for (uint i = req(); i < other->req(); i++) {
4137             assert(other->is_empty_memory(other->in(i)), "slice left uncovered");
4138         }
4139         #endif
4140     }
4141     // Replace spurious copies of base_memory by top.
4142     Node* base_mem = base_memory();
4143     if (base_mem != NULL && !base_mem->is_top()) {
4144         for (uint i = Compile::AliasIdxBot+1, imax = req(); i < imax; i++) {
4145             if (in(i) == base_mem)
4146                 set_req(i, empty_memory());
4147         }
4148     }
4149 }

4151 //-----grow_to_match-----
4152 void MergeMemNode::grow_to_match(const MergeMemNode* other) {
4153     Node* empty_mem = empty_memory();
4154     assert(other->is_empty_memory(empty_mem), "consistent sentinels");
4155     // look for the finite support of the other memory
4156     for (uint i = other->req(); --i >= req(); ) {
4157         if (other->in(i) != empty_mem) {
4158             uint new_len = i+1;
4159             while (req() < new_len) add_req(empty_mem);
4160             break;
4161         }
4162     }
4163 }

4165 //-----verify_sparse-----
4166 #ifndef PRODUCT
4167 bool MergeMemNode::verify_sparse() const {
4168     assert(is_empty_memory(make_empty_memory()), "sane sentinel");
4169     Node* base_mem = base_memory();
4170     // The following can happen in degenerate cases, since empty==top.
4171     if (is_empty_memory(base_mem)) return true;
4172     for (uint i = Compile::AliasIdxRaw; i < req(); i++) {
4173         assert(in(i) != NULL, "sane slice");
4174         if (in(i) == base_mem) return false; // should have been the sentinel valu
4175     }
4176     return true;
4177 }

4179 bool MergeMemStream::match_memory(Node* mem, const MergeMemNode* mm, int idx) {
4180     Node* n;

```

```

4181     n = mm->in(idx);
4182     if (mem == n) return true; // might be empty_memory()
4183     n = (idx == Compile::AliasIdxBot)? mm->base_memory(): mm->memory_at(idx);
4184     if (mem == n) return true;
4185     while (n->is_Phi() && (n = n->as_Phi()->is_copy()) != NULL) {
4186         if (mem == n) return true;
4187         if (n == NULL) break;
4188     }
4189     return false;
4190 }
4191 #endif // !PRODUCT

```

```

*****
20541 Thu Sep 1 02:19:28 2011
new/src/share/vm/opto/parse3.cpp
*****
_____unchanged_portion_omitted_____

147 void Parse::do_get_xxx(Node* obj, ciField* field, bool is_field) {
148     // Does this field have a constant value? If so, just push the value.
149     if (field->is_constant()) {
150         // final field
151     #endif /* !codereview */
152         if (field->is_static()) {
153             // final static field
154             if (push_constant(field->constant_value()))
155                 return;
156         }
157         else {
158             // final non-static field
159             // Treat final non-static fields of trusted classes (classes in
160             // java.lang.invoke and sun.invoke packages and subpackages) as
161             // compile time constants.
162             // final non-static field of a trusted package class (classes in
163             // java.lang.invoke and sun.invoke packages and subpackages).
164             if (obj->is_Con()) {
165                 const TypeOopPtr* oop_ptr = obj->bottom_type()->isa_oooptr();
166                 ciObject* constant_ooop = oop_ptr->const_ooop();
167                 ciConstant constant = field->constant_value_of(constant_ooop);
168             }
169             if (push_constant(constant, true))
170                 return;
171         }
172     }
173     ciType* field_klass = field->type();
174     bool is_vol = field->is_volatile();

175     // Compute address and memory type.
176     int offset = field->offset_in_bytes();
177     const TypePtr* adr_type = C->alias_type(field)->adr_type();
178     Node *adr = basic_plus_adr(obj, obj, offset);
179     BasicType bt = field->layout_type();

181     // Build the resultant type of the load
182     const Type *type;

184     bool must_assert_null = false;

186     if( bt == T_OBJECT ) {
187         if (!field->type()->is_loaded()) {
188             type = TypeInstPtr::BOTTOM;
189             must_assert_null = true;
190         } else if (field->is_constant() && field->is_static()) {
191             // This can happen if the constant oop is non-perm.
192             ciObject* con = field->constant_value().as_object();
193             // Do not "join" in the previous type; it doesn't add value,
194             // and may yield a vacuous result if the field is of interface type.
195             type = TypeOopPtr::make_from_constant(con)->isa_oooptr();
196             assert(type != NULL, "field singleton type must be consistent");
197         } else {
198             type = TypeOopPtr::make_from_klass(field_klass->as_klass());
199         }
200     } else {
201         type = Type::get_const_basic_type(bt);
202     }

```

```

203     // Build the load.
204     Node* ld = make_load(NULL, adr, type, bt, adr_type, is_vol);

206     // Adjust Java stack
207     if (type2size[bt] == 1)
208         push(ld);
209     else
210         push_pair(ld);

212     if (must_assert_null) {
213         // Do not take a trap here. It's possible that the program
214         // will never load the field's class, and will happily see
215         // null values in this field forever. Don't stumble into a
216         // trap for such a program, or we might get a long series
217         // of useless recompilations. (Or, we might load a class
218         // which should not be loaded.) If we ever see a non-null
219         // value, we will then trap and recompile. (The trap will
220         // not need to mention the class index, since the class will
221         // already have been loaded if we ever see a non-null value.)
222         // uncommon_trap(iter().get_field_signature_index());
223     #ifndef PRODUCT
224         if (PrintOpto && (Verbose || WizardMode)) {
225             method()->print_name(); tty->print_cr(" asserting nullness of field at bci
226         }
227     #endif
228     if (C->log() != NULL) {
229         C->log()->elem("assert_null reason='field' klass='%d'",
230             C->log()->identify(field->type()));
231     }
232     // If there is going to be a trap, put it at the next bytecode:
233     set_bci(iter().next_bci());
234     do_null_assert(peek(), T_OBJECT);
235     set_bci(iter().cur_bci()); // put it back
236 }

238     // If reference is volatile, prevent following memory ops from
239     // floating up past the volatile read. Also prevents commoning
240     // another volatile read.
241     if (field->is_volatile()) {
242         // Memory barrier includes bogus read of value to force load BEFORE membar
243         insert_mem_bar(Op_MemBarAcquire, ld);
244     }
245 }
_____unchanged_portion_omitted_____

```